

The LIFT Project

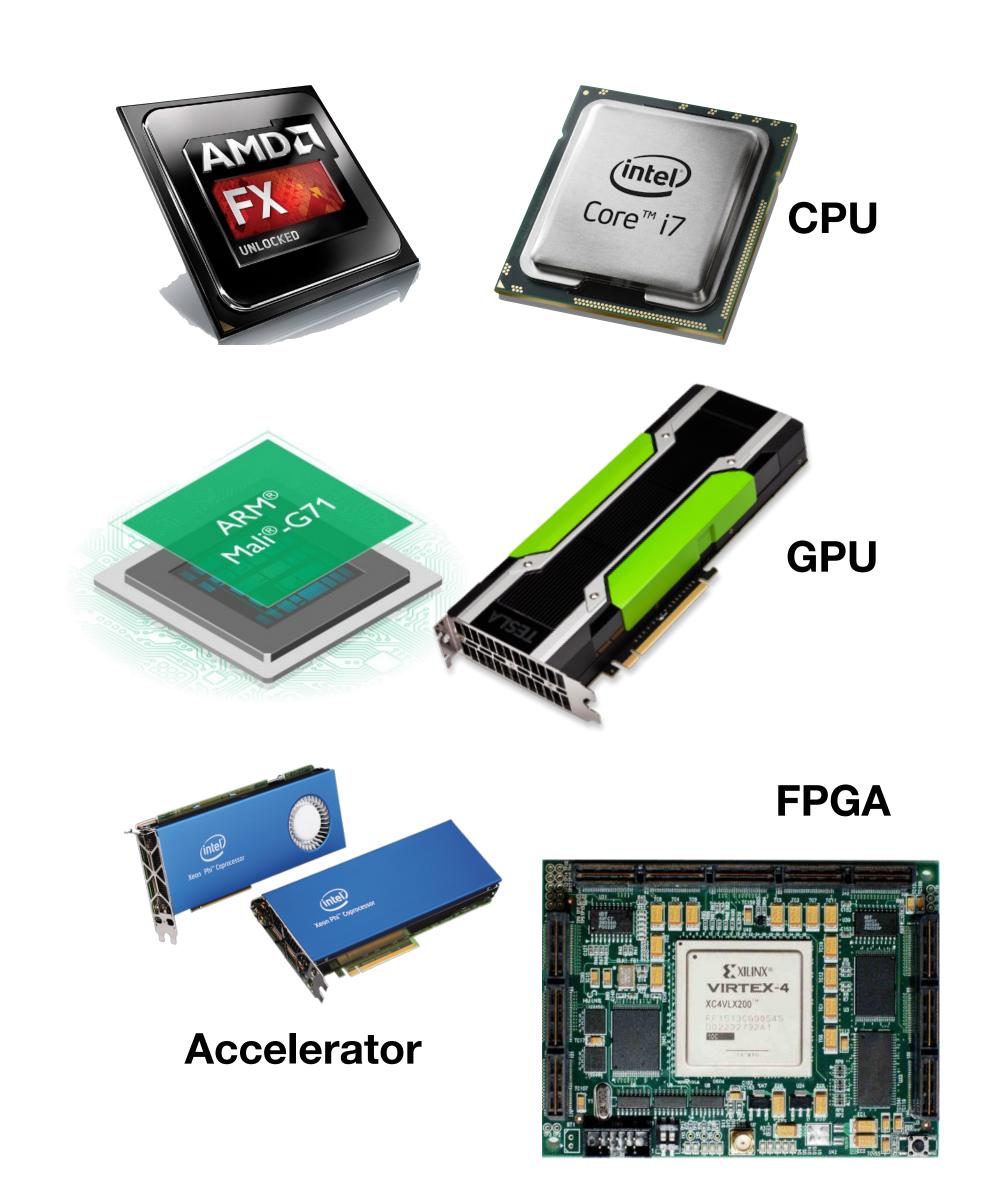
Performance Portable Parallel Code Generation via Rewrite Rules

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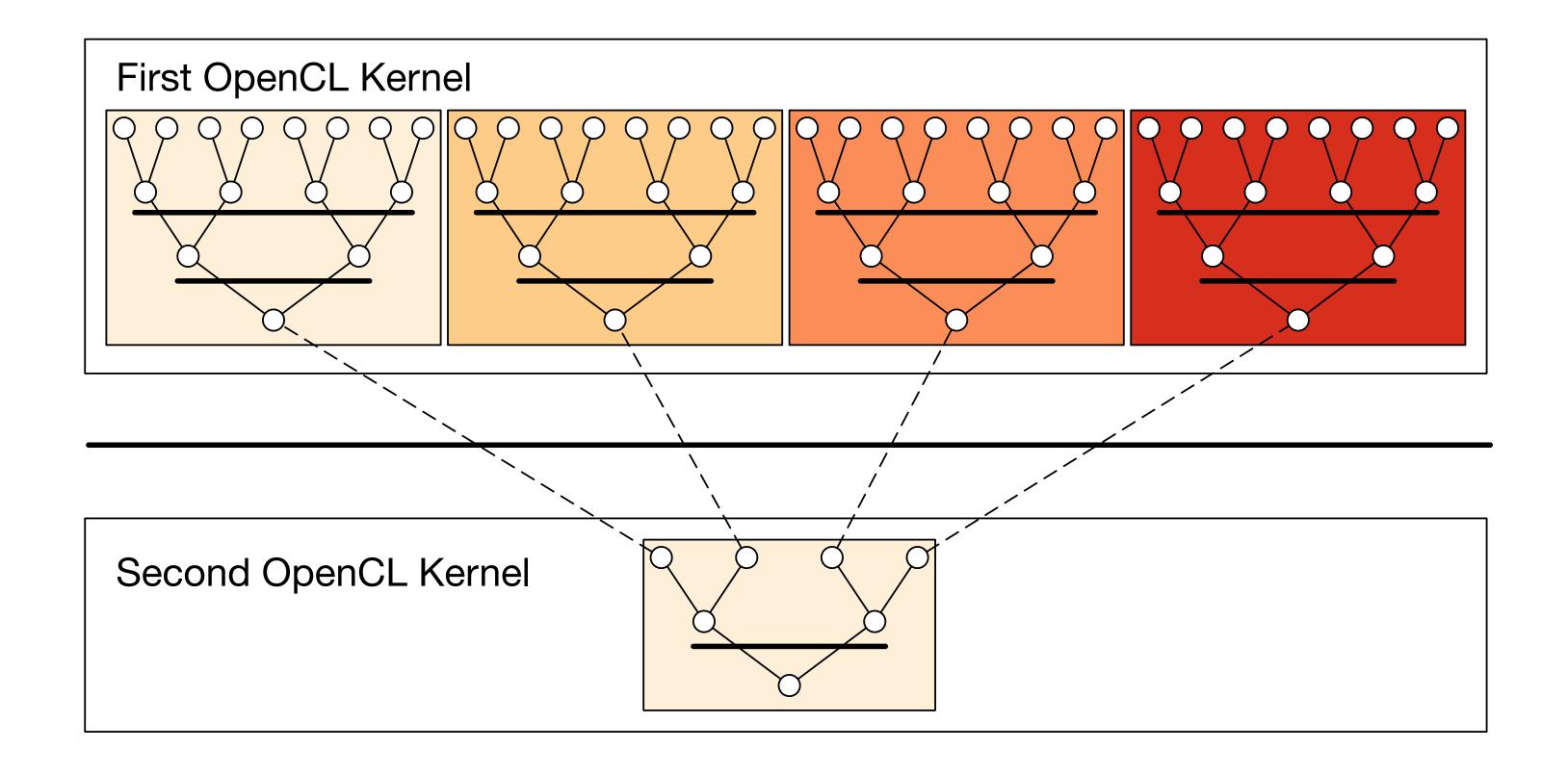
What are the problems LIFT tries to tackle?

- Parallel processors everywhere
- Many different types: CPUs, GPUs, ...
- Parallel programming is hard
- Optimising is even harder
- Problem:No portability of performance!



Case Study: Parallel Reduction in OpenCL

- Summing up all values of an array
- Comparison of 7 implementations by Nvidia
- Investigating complexity and efficiency of optimisations



Kernel function executed in parallel by multiple work-items

Work-items are identified by a unique global id

Work-items are grouped into work-groups

Local id within work-group

Small, but fast **local** memory

Big, but slow global memory

Memory **barriers** for consistency

Potential **Deadlock!**

Functionally correct implementations in OpenCL are hard!

1. Version: Unoptimised Implementation Parallel Reduction

```
kernel void reduce0(global float* g_idata, global float* g_odata,
                    unsigned int n, local float* l data) {
 unsigned int tid = get_local_id(0);
 unsigned int i = get_global_id(0);
  l_data[tid] = (i < n) ? g_idata[i] : 0;
  barrier(CLK_LOCAL_MEM_FENCE);
  // do reduction in local memory
 for (unsigned int s=1; s < get_local_size(0); s*= 2) {</pre>
   if ((tid % (2*s)) == 0) {
      l data[tid] += l_data[tid + s];
    barrier(CLK_LOCAL_MEM_FENCE);
  // write result for this work-group to global memory
  if (tid == 0) g_odata[get_group_id(0)] = l_data[0];
```

2. Version: Avoid Divergent Branching

```
kernel void reduce1(global float* g_idata, global float* g_odata,
                    unsigned int n, local float* l_data) {
 unsigned int tid = get_local_id(0);
 unsigned int i = get_global_id(0);
  l_data[tid] = (i < n) ? g_idata[i] : 0;</pre>
  barrier(CLK_LOCAL_MEM_FENCE);
 for (unsigned int s=1; s < get_local_size(0); s*= 2) {</pre>
    // continuous work-items remain active
    int index = 2 * s * tid;
    if (index < get_local_size(0)) {</pre>
      l data[index] += l data[index + s];
    barrier(CLK_LOCAL_MEM_FENCE);
 if (tid == 0) g_odata[get_group_id(0)] = l_data[0];
```

3. Version: Avoid Interleaved Addressing

```
kernel void reduce2(global float* g_idata, global float* g_odata,
                    unsigned int n, local float* l data) {
 unsigned int tid = get_local_id(0);
 unsigned int i = get_global_id(0);
  l_data[tid] = (i < n) ? g_idata[i] : 0;</pre>
  barrier(CLK_LOCAL_MEM_FENCE);
 // process elements in different order
 // requires commutativity
  for (unsigned int s=get_local_size(0)/2; s>0; s>>=1) {
   if (tid < s) {
      l_data[tid] += l_data[tid + s];
   barrier(CLK_LOCAL_MEM_FENCE);
 if (tid == 0) g_odata[get_group_id(0)] = l_data[0];
```

4. Version: Increase Computational Intensity per Work-Item

```
kernel void reduce3(global float* g_idata, global float* g_odata,
                    unsigned int n, local float* l data) {
  unsigned int tid = get_local_id(0);
  unsigned int i = get_group_id(0) * (get_local_size(0)*2)
                                    + get_local_id(0);
  l_data[tid] = (i < n) ? g_idata[i] : 0;
  // performs first addition during loading
  if (i + get_local_size(0) < n)</pre>
   l_data[tid] += g_idata[i+get_local_size(0)];
  barrier(CLK_LOCAL_MEM_FENCE);
  for (unsigned int s=get_local_size(0)/2; s>0; s>>=1) {
    if (tid < s) {
      l_data[tid] += l_data[tid + s];
    barrier(CLK_LOCAL_MEM_FENCE);
  if (tid == 0) g_odata[get_group_id(0)] = l_data[0];
```

5. Version: Avoid Synchronisation inside a Warp

```
kernel void reduce4(global float* g_idata, global float* g_odata,
                    unsigned int n, local volatile float* l_data) {
  unsigned int tid = get_local_id(0);
  unsigned int i = get_group_id(0) * (get_local_size(0)*2)
                                   + get_local_id(0);
  l_data[tid] = (i < n) ? g_idata[i] : 0;
  if (i + get_local_size(0) < n)</pre>
    l_data[tid] += g_idata[i+get_local_size(0)];
  barrier(CLK_LOCAL_MEM_FENCE);
  # pragma unroll 1
  for (unsigned int s=get_local_size(0)/2; s>32; s>>=1) {
    if (tid < s) { l data[tid] += l data[tid + s]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
  // this is not portable OpenCL code!
  if (tid < 32) {
    if (WG_SIZE >= 64) { l_data[tid] += l_data[tid+32]; }
    if (WG_SIZE >= 32) { l_data[tid] += l_data[tid+16]; }
    if (WG_SIZE >= 16) { l_data[tid] += l_data[tid+ 8]; }
    if (WG_SIZE >= 8) { l_data[tid] += l_data[tid+ 4]; }
    if (WG_SIZE >= 4) { l_data[tid] += l_data[tid+ 2]; }
    if (WG_SIZE >= 2) { l_data[tid] += l_data[tid+ 1]; } }
  if (tid == 0) g_odata[get_group_id(0)] = l_data[0]; }
```

6. Version: Complete Loop Unrolling

```
kernel void reduce5(global float* g_idata, global float* g_odata,
                    unsigned int n, local volatile float* l_data) {
  unsigned int tid = get_local_id(0);
  unsigned int i = get_group_id(0) * (get_local_size(0)*2)
                                   + get_local_id(0);
  l_data[tid] = (i < n) ? g_idata[i] : 0;
  if (i + get_local_size(0) < n)</pre>
    l_data[tid] += g_idata[i+get_local_size(0)];
  barrier(CLK_LOCAL_MEM_FENCE);
  if (WG_SIZE >= 256) {
    if (tid < 128) { l_data[tid] += l_data[tid+128]; }</pre>
    barrier(CLK_LOCAL_MEM_FENCE); }
  if (WG_SIZE >= 128) {
    if (tid < 64) { l_data[tid] += l_data[tid+ 64]; }</pre>
    barrier(CLK_LOCAL_MEM_FENCE); }
  if (tid < 32) {
    if (WG_SIZE >= 64) { l_data[tid] += l_data[tid+32]; }
    if (WG_SIZE >= 32) { l_data[tid] += l_data[tid+16]; }
    if (WG_SIZE >= 16) { l_data[tid] += l_data[tid+ 8]; }
    if (WG_SIZE >= 8) { l_data[tid] += l_data[tid+ 4]; }
    if (WG_SIZE >= 4) { l_data[tid] += l_data[tid+ 2]; }
    if (WG_SIZE >= 2) { l_data[tid] += l_data[tid+ 1]; } }
  if (tid == 0) g_odata[get_group_id(0)] = l_data[0]; }
```

7. Version: Fully Optimised Implementation

```
kernel void reduce6(global float* g_idata, global float* g_odata,
                    unsigned int n, local volatile float* l_data) {
  unsigned int tid = get_local_id(0);
  unsigned int i = get_group_id(0) * (get_local_size(0)*2)
                                   + get_local_id(0);
  unsigned int gridSize = WG_SIZE * get_num_groups(0);
  l_data[tid] = 0;
  while (i < n) { l_data[tid] += g_idata[i];</pre>
                  if (i + WG_SIZE < n)
                    l_data[tid] += g_idata[i+WG_SIZE];
                  i += gridSize; }
  barrier(CLK_LOCAL_MEM_FENCE);
  if (WG SIZE >= 256) {
   if (tid < 128) { l_data[tid] += l_data[tid+128]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
  if (WG_SIZE >= 128) {
   if (tid < 64) { l_data[tid] += l_data[tid+ 64]; }</pre>
    barrier(CLK_LOCAL_MEM_FENCE); }
  if (tid < 32) {
   if (WG_SIZE >= 64) { l_data[tid] += l_data[tid+32]; }
   if (WG_SIZE >= 32) { l_data[tid] += l_data[tid+16]; }
    if (WG_SIZE >= 16) { l_data[tid] += l_data[tid+ 8]; }
    if (WG_SIZE >= 8) { l_data[tid] += l_data[tid+ 4]; }
    if (WG_SIZE >= 4) { l_data[tid] += l_data[tid+ 2]; }
    if (WG_SIZE >= 2) { l_data[tid] += l_data[tid+ 1]; } }
  if (tid == 0) g_odata[get_group_id(0)] = l_data[0]; }
```

Reduction Case Study Conclusions

- Optimising OpenCL is complex
- Understanding of target hardware required
- Program changes not obvious
- Is it worth it? ...

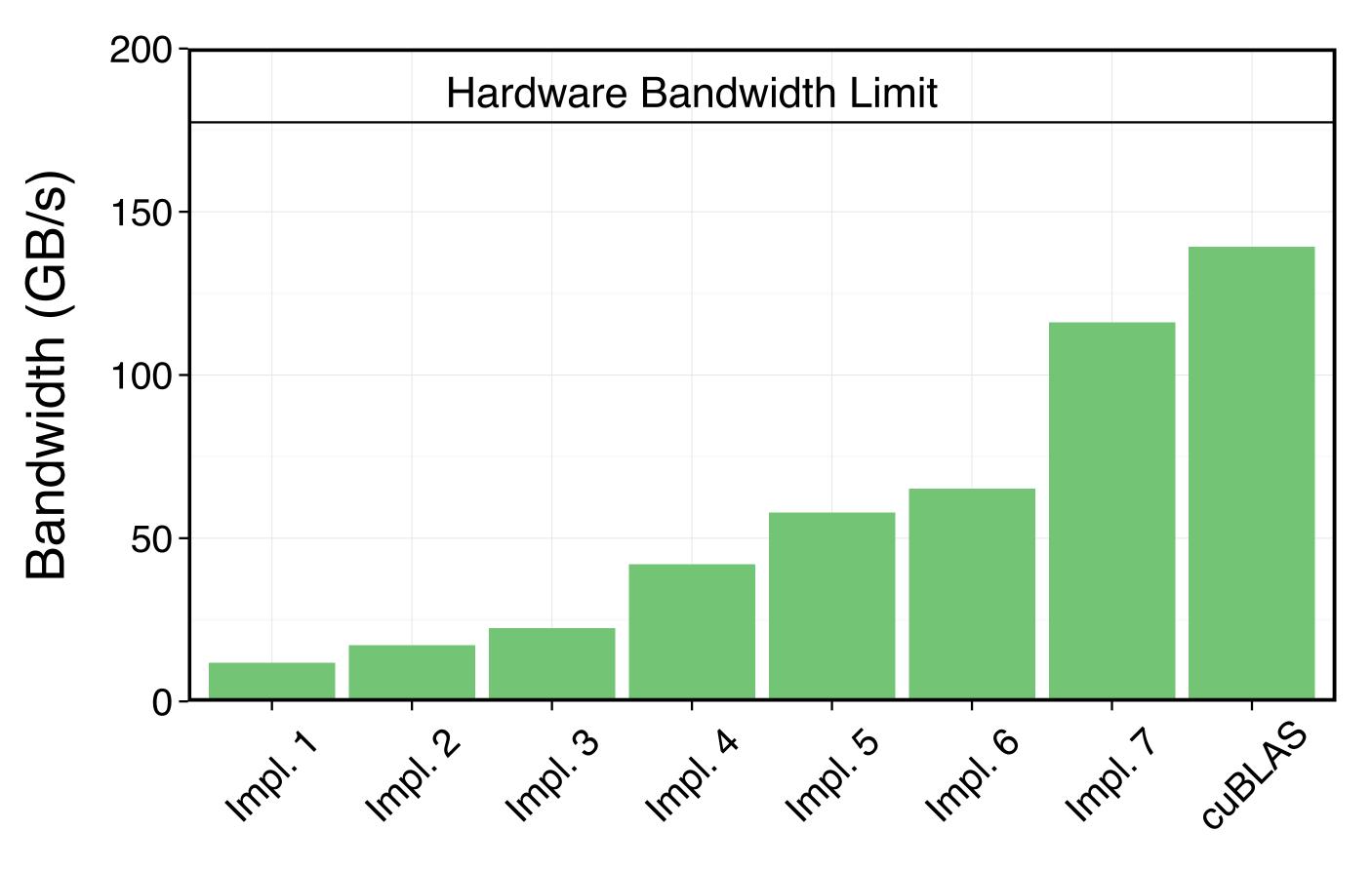
```
kernel
void reduce0(global float* g_idata,
             global float* g_odata,
             unsigned int n,
             local float* l_data) {
  unsigned int tid = get_local_id(0);
  unsigned int i = get_global_id(0);
  l_data[tid] = (i < n) ? g_idata[i] : 0;</pre>
  barrier(CLK_LOCAL_MEM_FENCE);
  for (unsigned int s=1;
       s < get_local_size(0); s*= 2) {</pre>
    if ((tid % (2*s)) == 0) {
      l data[tid] += l data[tid + s];
    barrier(CLK_LOCAL_MEM_FENCE);
  if (tid == 0)
    g_odata[get_group_id(0)] = l_data[0];
```

Unoptimized Implementation

```
kernel
void reduce6(global float* g_idata,
            global float* g_odata,
            unsigned int n,
            local volatile float* l_data) {
  unsigned int tid = get_local_id(0);
  unsigned int i =
   get_group_id(0) * (get_local_size(0)*2)
                   + get_local_id(0);
 unsigned int gridSize =
   WG_SIZE * get_num_groups(0);
 l_data[tid] = 0;
  while (i < n) {
   l_data[tid] += g_idata[i];
   if (i + WG_SIZE < n)
     l_data[tid] += g_idata[i+WG_SIZE];
   i += gridSize; }
 barrier(CLK_LOCAL_MEM_FENCE);
 if (WG_SIZE >= 256) {
   if (tid < 128) {
     l_data[tid] += l_data[tid+128]; }
   barrier(CLK_LOCAL_MEM_FENCE); }
 if (WG_SIZE >= 128) {
   if (tid < 64) {
     l_data[tid] += l_data[tid+ 64]; }
   barrier(CLK_LOCAL_MEM_FENCE); }
 if (tid < 32) {
   if (WG_SIZE >= 64) {
     l_data[tid] += l_data[tid+32]; }
   if (WG_SIZE >= 32) {
     l_data[tid] += l_data[tid+16]; }
   if (WG_SIZE >= 16) {
     l_data[tid] += l_data[tid+ 8]; }
   if (WG_SIZE >= 8) {
     l_data[tid] += l_data[tid+ 4]; }
   if (WG_SIZE >= 4) {
     l_data[tid] += l_data[tid+ 2]; }
   if (WG_SIZE >= 2) {
     l_data[tid] += l_data[tid+ 1]; } }
   g_odata[get_group_id(0)] = l_data[0];
```

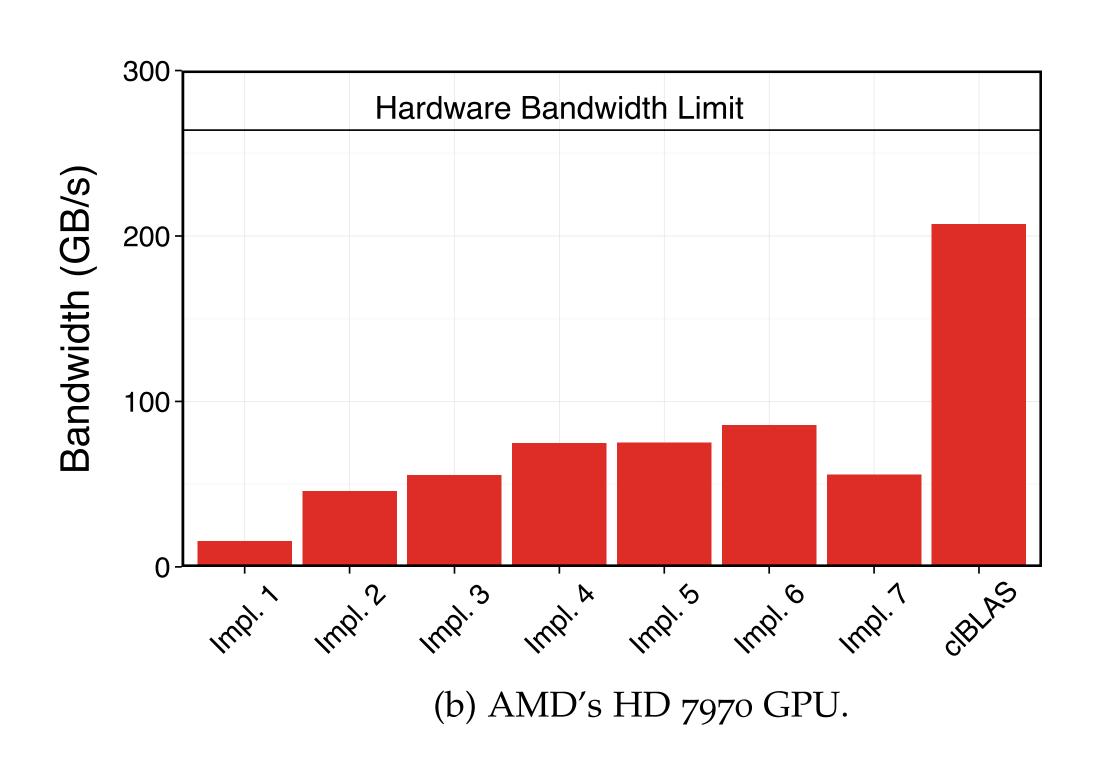
Fully Optimized Implementation

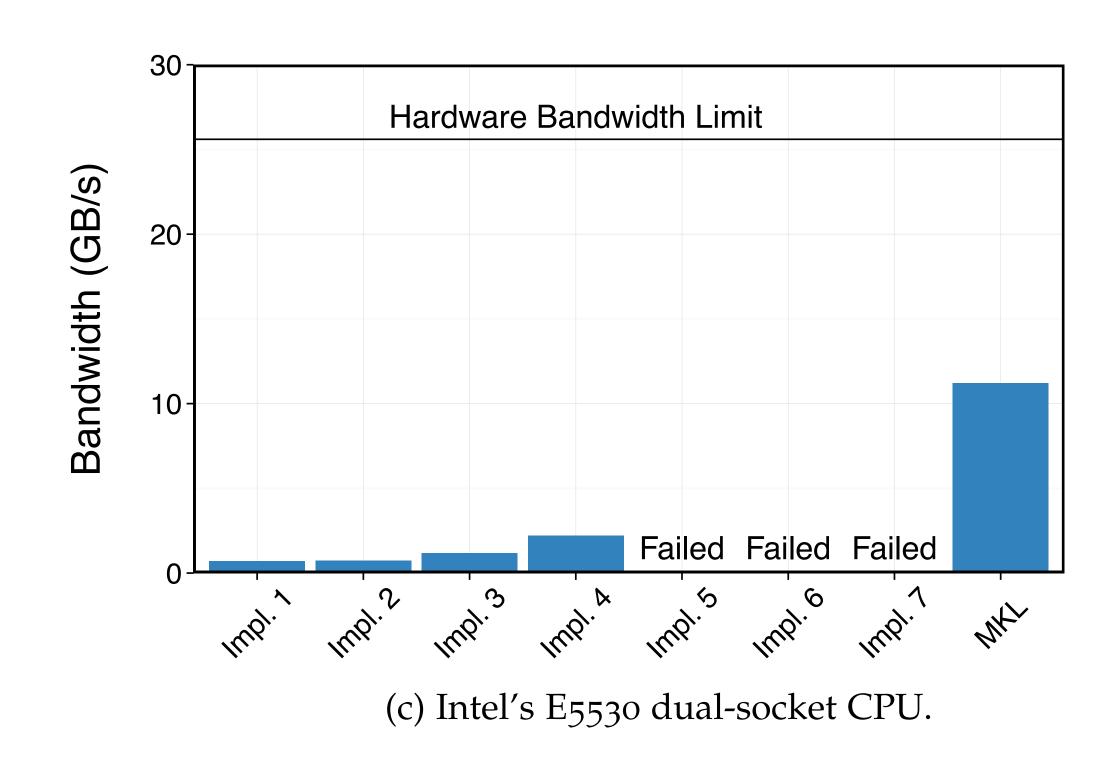
Performance Results Nvidia



- (a) Nvidia's GTX 480 GPU.
- ... Yes! Optimising improves performance by a factor of 10!
- Optimising is important, but ...

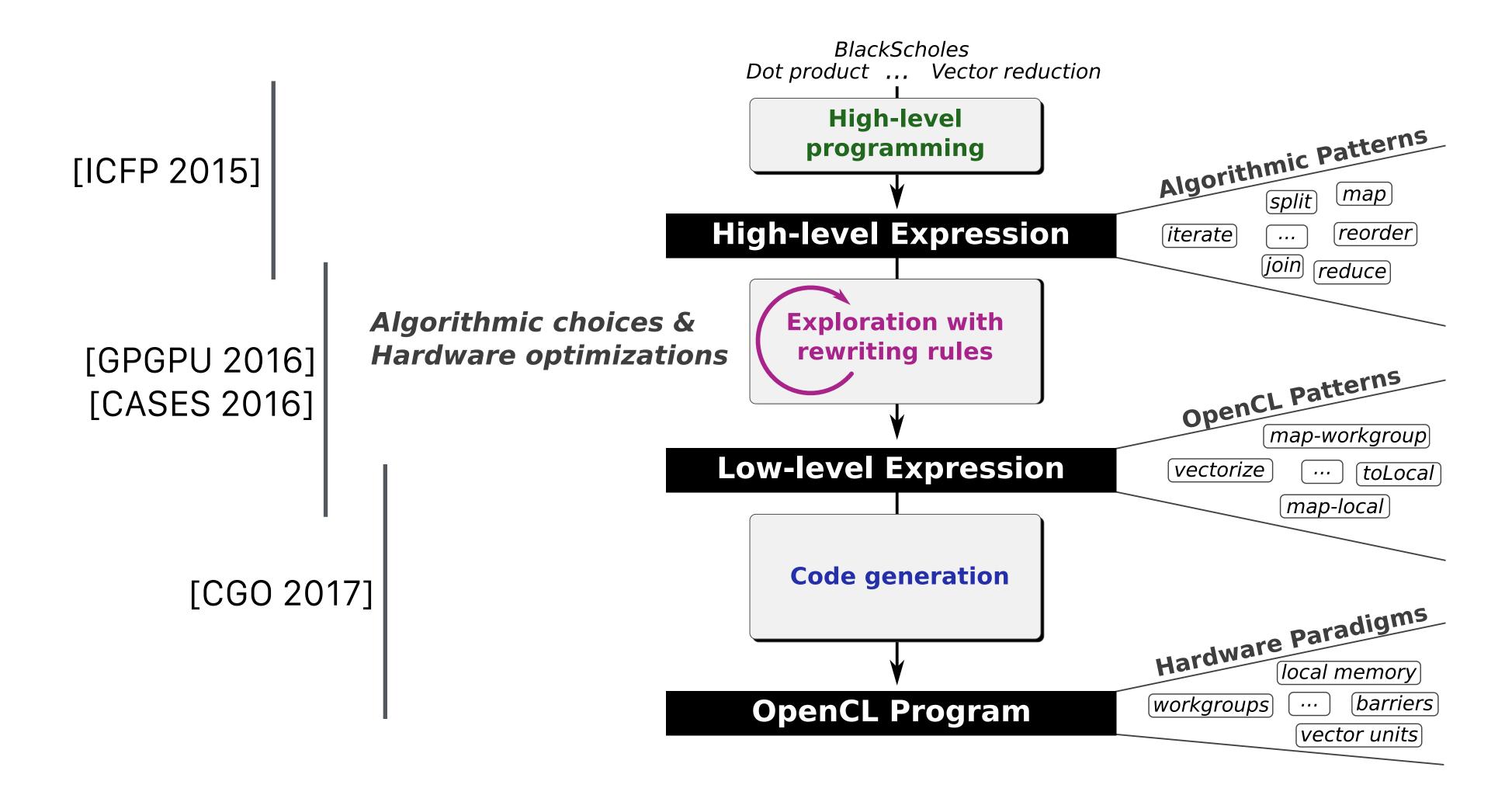
Performance Results AMD and Intel





- ... unfortunately, optimisations in OpenCL are not portable!
- Challenge: how to achieving portable performance?

LIFT: Performance Portable GPU Code Generation via Rewrite Rules



Ambition: automatic generation of Performance Portable code

Walkthrough

```
1 sum(vec) = reduce(+,0,vec)

rewrite rules code generation
```

```
vecSum = reduce ∘ join ∘ map-workgroup (
    join ∘ toGlobal (map-local (map-seq id)) ∘ split 1 ∘
    join ∘ map-warp (
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 1 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 2 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 4 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 8 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 16 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 32
    ) ∘ split 64 ∘
    join ∘ map-local (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 64 ∘
    join ∘ toLocal (map-local (reduce-seq (+) 0)) ∘
    split (blockSize/128) ∘ reorder-stride 128
) ∘ split blockSize
```

```
kernel
void reduce6(global float* g_idata,
             global float* g_odata,
             unsigned int n,
             local volatile float* l_data) {
 unsigned int tid = get_local_id(0);
 unsigned int i =
   get_group_id(0) * (get_local_size(0)*2)
                    + get_local_id(0);
 unsigned int gridSize =
   WG_SIZE * get_num_groups(0);
 l_data[tid] = 0;
 while (i < n) {
   l_data[tid] += g_idata[i];
   if (i + WG_SIZE < n)</pre>
      l_data[tid] += g_idata[i+WG_SIZE];
   i += gridSize; }
 barrier(CLK_LOCAL_MEM_FENCE);
 if (WG_SIZE >= 256) {
   if (tid < 128) {
      l_data[tid] += l_data[tid+128]; }
   barrier(CLK_LOCAL_MEM_FENCE); }
 if (WG_SIZE >= 128) {
   if (tid < 64) {
      l_data[tid] += l_data[tid+ 64]; }
   barrier(CLK_LOCAL_MEM_FENCE); }
 if (tid < 32) {
   if (WG_SIZE >= 64) {
      l_data[tid] += l_data[tid+32]; }
   if (WG_SIZE >= 32) {
     l_data[tid] += l_data[tid+16]; }
   if (WG_SIZE >= 16) {
     l_data[tid] += l_data[tid+ 8]; }
   if (WG_SIZE >= 8) {
     l_data[tid] += l_data[tid+ 4]; }
   if (WG_SIZE >= 4) {
     l_data[tid] += l_data[tid+ 2]; }
   if (WG_SIZE >= 2) {
     l_data[tid] += l_data[tid+ 1]; } }
 if (tid == 0)
   g_odata[get_group_id(0)] = l_data[0];
```

Walkthrough

 $1) \quad sum(vec) = reduce(+, 0, vec)$

rewrite rules

code generation



```
vecSum = reduce o join o map-workgroup (
    join o toGlobal (map-local (map-seq id)) o split 1 o
    join o map-warp (
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 1 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 2 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 4 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 8 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 16 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 32
        ) o split 64 o
        join o map-local (reduce-seq (+) 0) o split 2 o reorder-stride 64 o
        join o toLocal (map-local (reduce-seq (+) 0)) o
        split (blockSize/128) o reorder-stride 128
        ) o split blockSize
```

```
kernel
void reduce6(global float* g_idata,
             global float* g_odata,
             unsigned int n,
             local volatile float* l_data) {
  unsigned int tid = get_local_id(0);
  unsigned int i =
    get_group_id(0) * (get_local_size(0)*2)
                    + get_local_id(0);
  unsigned int gridSize =
   WG_SIZE * get_num_groups(0);
  l_data[tid] = 0;
  while (i < n) {
    l_data[tid] += g_idata[i];
    if (i + WG_SIZE < n)</pre>
      l_data[tid] += g_idata[i+WG_SIZE];
    i += gridSize; }
  barrier(CLK_LOCAL_MEM_FENCE);
  if (WG_SIZE >= 256) {
    if (tid < 128) {
      l_data[tid] += l_data[tid+128]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
  if (WG_SIZE >= 128) {
    if (tid < 64) {
     l_data[tid] += l_data[tid+ 64]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
  if (tid < 32) {
    if (WG_SIZE >= 64) {
      l_data[tid] += l_data[tid+32]; }
    if (WG_SIZE >= 32) {
      l_data[tid] += l_data[tid+16]; }
    if (WG_SIZE >= 16) {
      l_data[tid] += l_data[tid+ 8]; }
    if (WG_SIZE >= 8) {
     l_data[tid] += l_data[tid+ 4]; }
    if (WG_SIZE >= 4) {
     l_data[tid] += l_data[tid+ 2]; }
    if (WG_SIZE >= 2) {
     l_data[tid] += l_data[tid+ 1]; } }
  if (tid == 0)
    g_odata[get_group_id(0)] = l_data[0];
```

1 Algorithmic Primitives (a.k.a. algorithmic skeletons)

map(f, x):	$\begin{bmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 \end{bmatrix}$	
zip(x, y):	$\begin{bmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 \\ y_1 & y_2 & y_3 & y_4 & y_5 & y_6 & y_7 & y_8 \end{bmatrix}$	$(x_{1}, y_{1})(x_{2}, y_{2})(x_{3}, y_{3})(x_{4}, y_{4})(x_{5}, y_{5})(x_{6}, y_{6})(x_{7}, y_{7})(x_{8}, y_{8})$
reduce(+, 0, x):	$\begin{bmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 \end{bmatrix}$	$\longrightarrow x_1 + x_2 + x_3 + x_4 + x_5 + x_6 + x_7 + x_8$
split(n, x):	$\begin{bmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 \end{bmatrix}$	
join(x):	$\begin{bmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 \end{bmatrix}$	$ \longrightarrow \begin{array}{ c c c c c c c c c c c c c c c c c c c$
iterate(f, n, x):	$\begin{bmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 \end{bmatrix}$	$\longmapsto f(\dots f(\boxed{x_1} \boxed{x_2} \boxed{x_3} \boxed{x_4} \boxed{x_5} \boxed{x_6} \boxed{x_7} \boxed{x_8})\dots)$
reorder(σ , x):	$\begin{bmatrix} x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 & x_8 \end{bmatrix}$	$\longrightarrow x_{\sigma(1)} x_{\sigma(2)} x_{\sigma(3)} x_{\sigma(4)} x_{\sigma(5)} x_{\sigma(6)} x_{\sigma(7)} x_{\sigma(8)}$

1 High-Level Programs

```
scal(a, vec) = map(\lambda x \mapsto x*a, vec)
asum(vec) = reduce(+, 0, map(abs, vec))
dotProduct(x, y) = reduce(+, 0, map(*, zip(x, y)))
gemv(mat, x, y, \alpha, \beta) =
 map(+, zip(
    map(\lambda row \rightarrow scal(\alpha, dotProduct(row, x)), mat),
    scal(\beta, y))
```

Walkthrough

1 sum(vec) = reduce(+,0,vec)
rewrite rules code generation

```
vecSum = reduce ∘ join ∘ map-workgroup (
    join ∘ toGlobal (map-local (map-seq id)) ∘ split 1 ∘
    join ∘ map-warp (
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 1 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 2 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 4 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 8 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 16 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 32
    ) ∘ split 64 ∘
    join ∘ map-local (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 64 ∘
    join ∘ toLocal (map-local (reduce-seq (+) 0)) ∘
    split (blockSize/128) ∘ reorder-stride 128
) ∘ split blockSize
```

```
kernel
void reduce6(global float* g_idata,
             global float* g_odata,
             unsigned int n,
             local volatile float* l_data) {
 unsigned int tid = get_local_id(0);
 unsigned int i =
   get_group_id(0) * (get_local_size(0)*2)
                    + get_local_id(0);
 unsigned int gridSize =
   WG_SIZE * get_num_groups(0);
 l_data[tid] = 0;
 while (i < n) {
   l_data[tid] += g_idata[i];
   if (i + WG_SIZE < n)</pre>
      l_data[tid] += g_idata[i+WG_SIZE];
   i += gridSize; }
 barrier(CLK_LOCAL_MEM_FENCE);
 if (WG_SIZE >= 256) {
   if (tid < 128) {
      l_data[tid] += l_data[tid+128]; }
   barrier(CLK_LOCAL_MEM_FENCE); }
 if (WG_SIZE >= 128) {
   if (tid < 64) {
      l_data[tid] += l_data[tid+ 64]; }
   barrier(CLK_LOCAL_MEM_FENCE); }
 if (tid < 32) {
   if (WG_SIZE >= 64) {
      l_data[tid] += l_data[tid+32]; }
   if (WG_SIZE >= 32) {
     l_data[tid] += l_data[tid+16]; }
   if (WG_SIZE >= 16) {
     l_data[tid] += l_data[tid+ 8]; }
   if (WG_SIZE >= 8) {
     l_data[tid] += l_data[tid+ 4]; }
   if (WG_SIZE >= 4) {
     l_data[tid] += l_data[tid+ 2]; }
   if (WG_SIZE >= 2) {
     l_data[tid] += l_data[tid+ 1]; } }
 if (tid == 0)
   g_odata[get_group_id(0)] = l_data[0];
```

Walkthrough

1 sum(vec) = reduce(+, 0, vec)

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```
vecSum = reduce o join o map-workgroup (
    join o toGlobal (map-local (map-seq id)) o split 1 o
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        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 1 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 2 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 4 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 8 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 16 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 32
        ) o split 64 o
        join o map-local (reduce-seq (+) 0) o split 2 o reorder-stride 64 o
        join o toLocal (map-local (reduce-seq (+) 0)) o
        split (blockSize/128) o reorder-stride 128
        ) o split blockSize
```

```
kernel
void reduce6(global float* g_idata,
             global float* g_odata,
            unsigned int n,
            local volatile float* l_data) {
  unsigned int tid = get_local_id(0);
  unsigned int i =
    get_group_id(0) * (get_local_size(0)*2)
                    + get_local_id(0);
  unsigned int gridSize =
   WG_SIZE * get_num_groups(0);
  l_data[tid] = 0;
  while (i < n) {
   l_data[tid] += g_idata[i];
    if (i + WG_SIZE < n)</pre>
      l_data[tid] += g_idata[i+WG_SIZE];
    i += gridSize; }
  barrier(CLK_LOCAL_MEM_FENCE);
  if (WG_SIZE >= 256) {
    if (tid < 128) {
      l_data[tid] += l_data[tid+128]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
  if (WG_SIZE >= 128) {
    if (tid < 64) {
     l_data[tid] += l_data[tid+ 64]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
  if (tid < 32) {
    if (WG_SIZE >= 64) {
      l_data[tid] += l_data[tid+32]; }
    if (WG_SIZE >= 32) {
     l_data[tid] += l_data[tid+16]; }
    if (WG_SIZE >= 16) {
     l_data[tid] += l_data[tid+ 8]; }
    if (WG_SIZE >= 8) {
     l_data[tid] += l_data[tid+ 4]; }
    if (WG_SIZE >= 4) {
      l_data[tid] += l_data[tid+ 2]; }
    if (WG_SIZE >= 2) {
     l_data[tid] += l_data[tid+ 1]; } }
  if (tid == 0)
    g_odata[get_group_id(0)] = l_data[0];
```

2 Algorithmic Rewrite Rules

- Provably correct rewrite rules
- Express algorithmic implementation choices

Split-join rule:

```
map \ f \rightarrow join \circ map \ (map \ f) \circ split \ n
```

Map fusion rule:

```
map \ f \circ map \ g \rightarrow map \ (f \circ g)
```

Reduce rules:

```
reduce f \ z \to reduce \ f \ z \circ reduce Part \ f \ z

reduce Part \ f \ z \to reduce Part \ f \ z \circ reorder

reduce Part \ f \ z \to join \ \circ map \ (reduce Part \ f \ z) \circ split \ n

reduce Part \ f \ z \to iterate \ n \ (reduce Part \ f \ z)
```

2 OpenCL Primitives

OpenCL concept **Primitive**

mapGlobal

map Workgroup

mapLocal

mapSeq

reduceSeq

toLocal, toGlobal

map Vec, $split Vec.\ join\ Vec$ Work-items

Work-groups

OpenCL thread hierarchy

workgroups

local threads

global threads

Sequential implementations

Memory areas

Vectorisation

2 OpenCL Rewrite Rules

Express low-level implementation and optimisation choices

Map rules:

```
map \ f \rightarrow map \ Workgroup \ f \mid map Local \ f \mid map Global \ f \mid map Seq \ f
```

Local/ global memory rules:

```
mapLocal\ f \rightarrow toLocal\ (mapLocal\ f) mapLocal\ f \rightarrow toGlobal\ (mapLocal\ f)
```

Vectorisation rule:

```
map\ f \rightarrow join\ Vec \circ map\ (map\ Vec\ f) \circ split\ Vec\ n
```

Fusion rule:

$$reduceSeq\ f\ z\circ mapSeq\ g \to reduceSeq\ (\lambda\ (acc,x).\ f\ (acc,g\ x))\ z$$

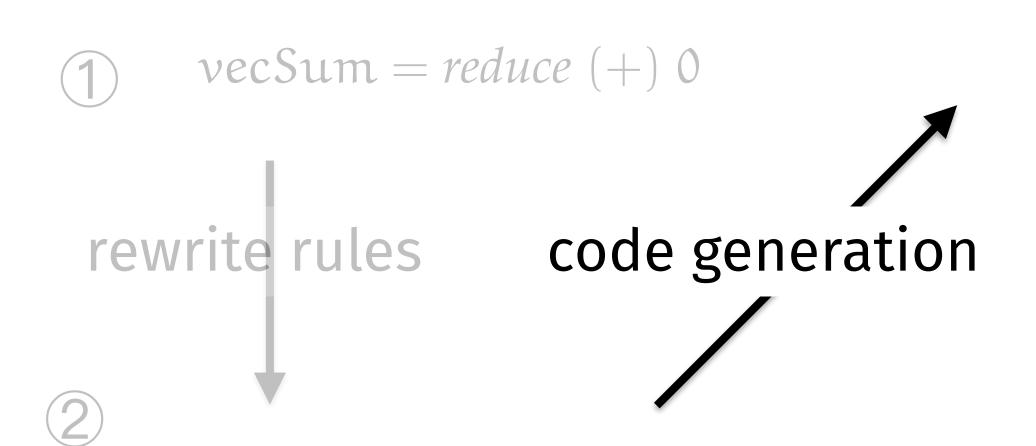
Walkthrough

1 vecSum = reduce (+) 0
rewrite rules code generation

```
vecSum = reduce ∘ join ∘ map-workgroup (
    join ∘ toGlobal (map-local (map-seq id)) ∘ split 1 ∘
    join ∘ map-warp (
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 1 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 2 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 4 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 8 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 16 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 32
    ) ∘ split 64 ∘
    join ∘ map-local (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 64 ∘
    join ∘ toLocal (map-local (reduce-seq (+) 0)) ∘
    split (blockSize/128) ∘ reorder-stride 128
) ∘ split blockSize
```

```
kernel
void reduce6(global float* g_idata,
             global float* g_odata,
             unsigned int n,
             local volatile float* l_data) {
 unsigned int tid = get_local_id(0);
 unsigned int i =
    get_group_id(0) * (get_local_size(0)*2)
                    + get_local_id(0);
 unsigned int gridSize =
    WG_SIZE * get_num_groups(0);
  l_data[tid] = 0;
 while (i < n) {
    l_data[tid] += g_idata[i];
    if (i + WG_SIZE < n)</pre>
      l_data[tid] += g_idata[i+WG_SIZE];
    i += gridSize; }
 barrier(CLK_LOCAL_MEM_FENCE);
 if (WG_SIZE >= 256) {
    if (tid < 128) {
      l_data[tid] += l_data[tid+128]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
 if (WG_SIZE >= 128) {
    if (tid < 64) {
      l_data[tid] += l_data[tid+ 64]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
 if (tid < 32) {
    if (WG_SIZE >= 64) {
      l_data[tid] += l_data[tid+32]; }
    if (WG_SIZE >= 32) {
      l_data[tid] += l_data[tid+16]; }
    if (WG_SIZE >= 16) {
      l_data[tid] += l_data[tid+ 8]; }
    if (WG_SIZE >= 8) {
      l_data[tid] += l_data[tid+ 4]; }
    if (WG_SIZE >= 4) {
      l_data[tid] += l_data[tid+ 2]; }
    if (WG_SIZE >= 2) {
     l_data[tid] += l_data[tid+ 1]; } }
 if (tid == 0)
    g_odata[get_group_id(0)] = l_data[0];
```

Walkthrough



```
vecSum = reduce o join o map-workgroup (
    join o toGlobal (map-local (map-seq id)) o split 1 o
    join o map-warp (
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 1 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 2 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 4 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 8 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 16 o
        join o map-lane (reduce-seq (+) 0) o split 2 o reorder-stride 32
      ) o split 64 o
        join o map-local (reduce-seq (+) 0) o split 2 o reorder-stride 64 o
        join o toLocal (map-local (reduce-seq (+) 0)) o
        split (blockSize/128) o reorder-stride 128
      ) o split blockSize
```

```
kernel
void reduce6(global float* g_idata,
             global float* g_odata,
             unsigned int n,
             local volatile float* l_data) {
 unsigned int tid = get_local_id(0);
 unsigned int i =
    get_group_id(0) * (get_local_size(0)*2)
                    + get_local_id(0);
 unsigned int gridSize =
    WG_SIZE * get_num_groups(0);
  l_data[tid] = 0;
 while (i < n) {
    l_data[tid] += g_idata[i];
    if (i + WG_SIZE < n)</pre>
      l_data[tid] += g_idata[i+WG_SIZE];
    i += gridSize; }
 barrier(CLK_LOCAL_MEM_FENCE);
 if (WG_SIZE >= 256) {
    if (tid < 128) {
      l_data[tid] += l_data[tid+128]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
 if (WG_SIZE >= 128) {
    if (tid < 64) {
      l_data[tid] += l_data[tid+ 64]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
 if (tid < 32) {
    if (WG_SIZE >= 64) {
      l_data[tid] += l_data[tid+32]; }
   if (WG_SIZE >= 32) {
      l_data[tid] += l_data[tid+16]; }
    if (WG_SIZE >= 16) {
      l_data[tid] += l_data[tid+ 8]; }
    if (WG_SIZE >= 8) {
      l_data[tid] += l_data[tid+ 4]; }
    if (WG_SIZE >= 4) {
      l_data[tid] += l_data[tid+ 2]; }
    if (WG_SIZE >= 2) {
     l_data[tid] += l_data[tid+ 1]; } }
 if (tid == 0)
   g_odata[get_group_id(0)] = l_data[0];
```

3 Pattern based OpenCL Code Generation

Generate OpenCL code for each OpenCL primitive

```
mapGlobal\ f\ xs \longrightarrow \begin{bmatrix} \textbf{for (int g_id = get_global_id(0); g_id < n; g_id += get_global_size(0)) } \\ \text{output[g_id] = } \textbf{f(xs[g_id]);} \\ \end{bmatrix}
```

```
reduceSeq f z xs \longrightarrow \begin{cases} \text{T acc} = z; \\ \text{for (int } i = 0; i < n; ++i) \\ \text{acc} = f(\text{acc, xs[i]}); \\ \end{cases}
```

•
•
•
•

 A lot more details about the code generation implementation can be found in our <u>CGO 2017 paper</u>

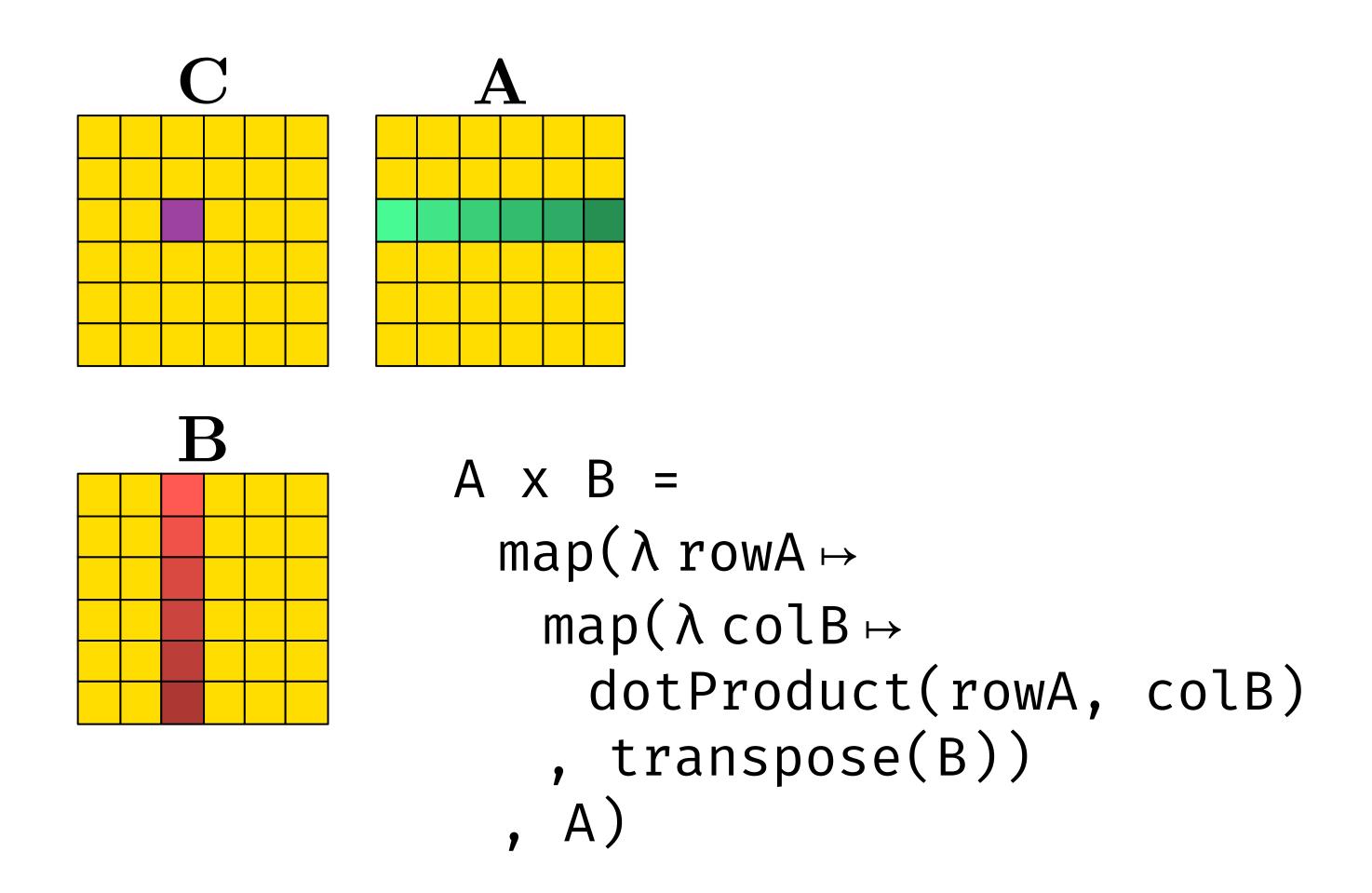
Walkthrough

1 vecSum = reduce (+) 0
rewrite rules code generation

```
vecSum = reduce ∘ join ∘ map-workgroup (
    join ∘ toGlobal (map-local (map-seq id)) ∘ split 1 ∘
    join ∘ map-warp (
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 1 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 2 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 4 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 8 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 16 ∘
        join ∘ map-lane (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 32
    ) ∘ split 64 ∘
    join ∘ map-local (reduce-seq (+) 0) ∘ split 2 ∘ reorder-stride 64 ∘
    join ∘ toLocal (map-local (reduce-seq (+) 0)) ∘
    split (blockSize/128) ∘ reorder-stride 128
) ∘ split blockSize
```

```
kernel
void reduce6(global float* g_idata,
             global float* g_odata,
             unsigned int n,
             local volatile float* l_data) {
 unsigned int tid = get_local_id(0);
 unsigned int i =
    get_group_id(0) * (get_local_size(0)*2)
                    + get_local_id(0);
 unsigned int gridSize =
    WG_SIZE * get_num_groups(0);
  l_data[tid] = 0;
 while (i < n) {
    l_data[tid] += g_idata[i];
    if (i + WG_SIZE < n)</pre>
      l_data[tid] += g_idata[i+WG_SIZE];
    i += gridSize; }
 barrier(CLK_LOCAL_MEM_FENCE);
 if (WG_SIZE >= 256) {
    if (tid < 128) {
      l_data[tid] += l_data[tid+128]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
 if (WG_SIZE >= 128) {
    if (tid < 64) {
      l_data[tid] += l_data[tid+ 64]; }
    barrier(CLK_LOCAL_MEM_FENCE); }
 if (tid < 32) {
    if (WG_SIZE >= 64) {
      l_data[tid] += l_data[tid+32]; }
    if (WG_SIZE >= 32) {
      l_data[tid] += l_data[tid+16]; }
    if (WG_SIZE >= 16) {
      l_data[tid] += l_data[tid+ 8]; }
    if (WG_SIZE >= 8) {
      l_data[tid] += l_data[tid+ 4]; }
    if (WG_SIZE >= 4) {
      l_data[tid] += l_data[tid+ 2]; }
    if (WG_SIZE >= 2) {
     l_data[tid] += l_data[tid+ 1]; } }
 if (tid == 0)
    g_odata[get_group_id(0)] = l_data[0];
```

Case Study: Matrix Multiplication



Tiling as a Rewrite Rules

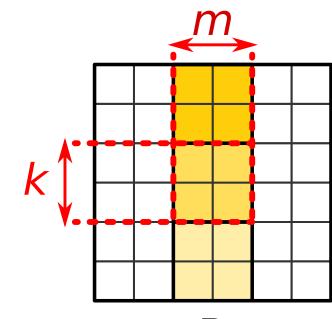
Naïve matrix multiplication

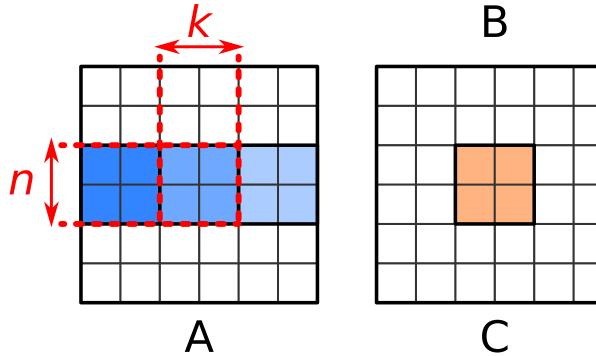
```
1 \quad map(\lambda \ arow \ .
2 \quad map(\lambda \ bcol \ .
3 \quad reduce(+, 0) \circ map(\times) \circ zip(arow, bcol)
4 \quad , transpose(B))
5 \quad , A)
```

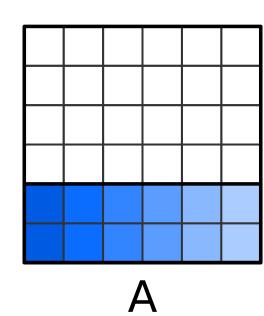


Apply tiling rules

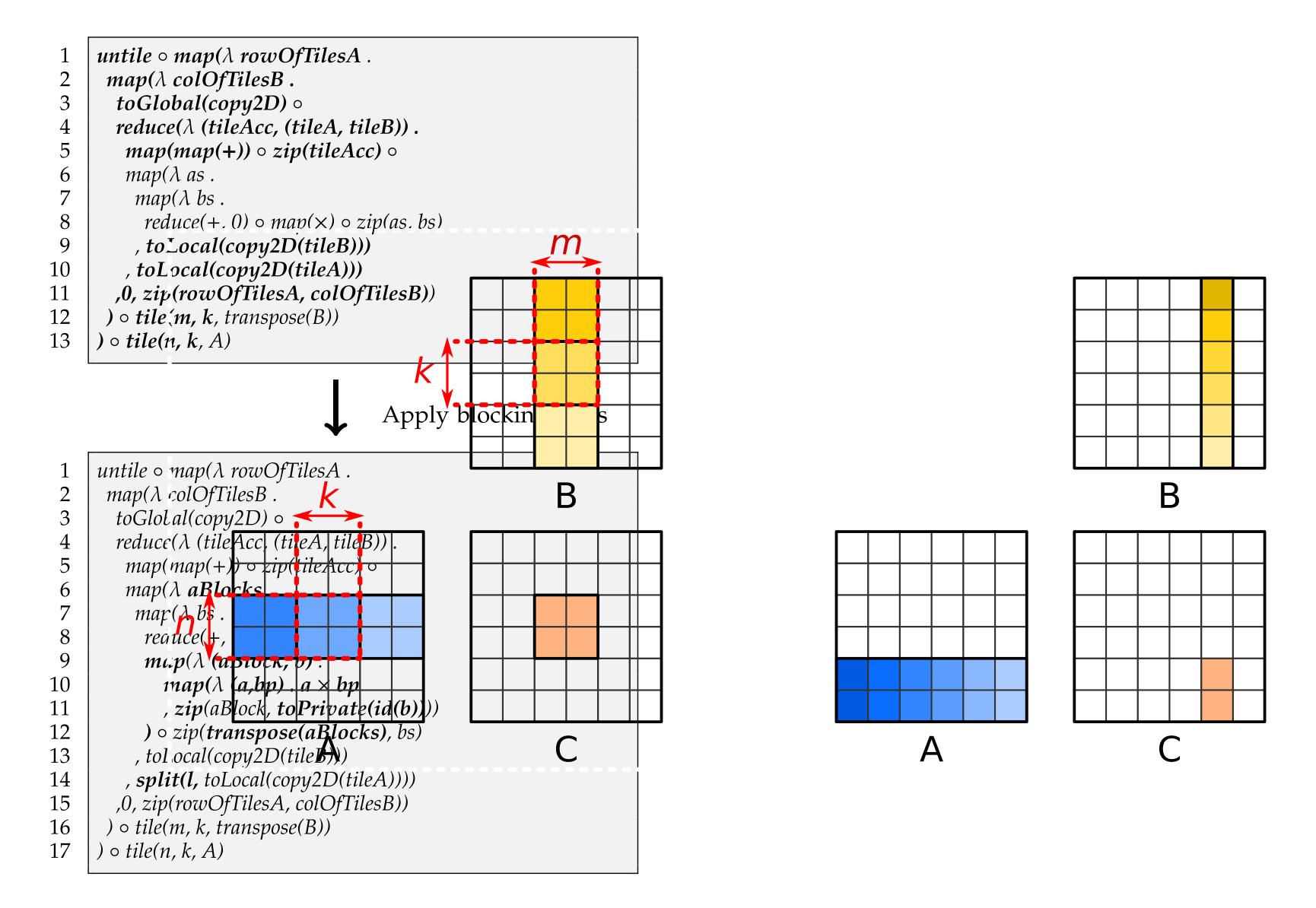
```
untile \circ map(\lambda rowOfTilesA.
       map(\lambda \ colOfTilesB.
 3
        toGlobal(copy2D) \circ
        reduce(\lambda (tileAcc, (tileA, tileB)).
         map(map(+)) \circ zip(tileAcc) \circ
         map(\lambda as.
 6
          map(\lambda bs.
           reduce(+, 0) \circ map(\times) \circ zip(as, bs)
 9
          , toLocal(copy2D(tileB)))
         , toLocal(copy2D(tileA)))
10
        ,0, zip(rowOfTilesA, colOfTilesB))
       ) \circ tile(m, k, transpose(B))
      ) \circ tile(n, k, A)
```





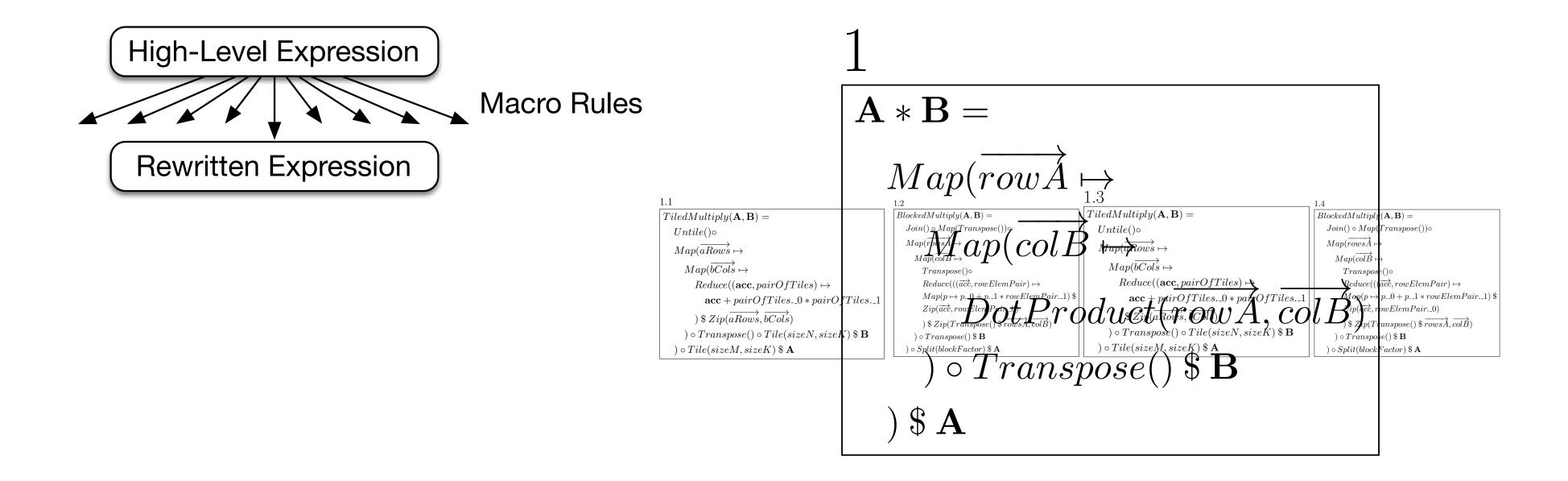


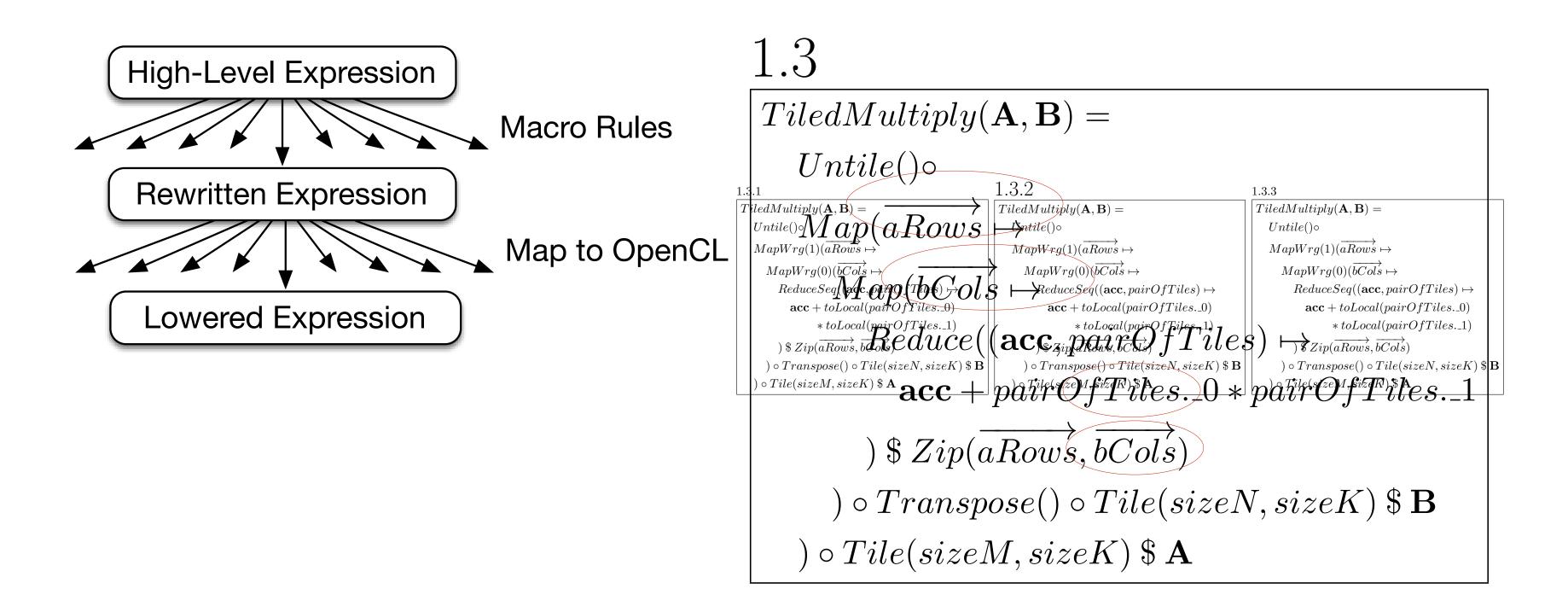
Register Blocking as a Rewrite Rules

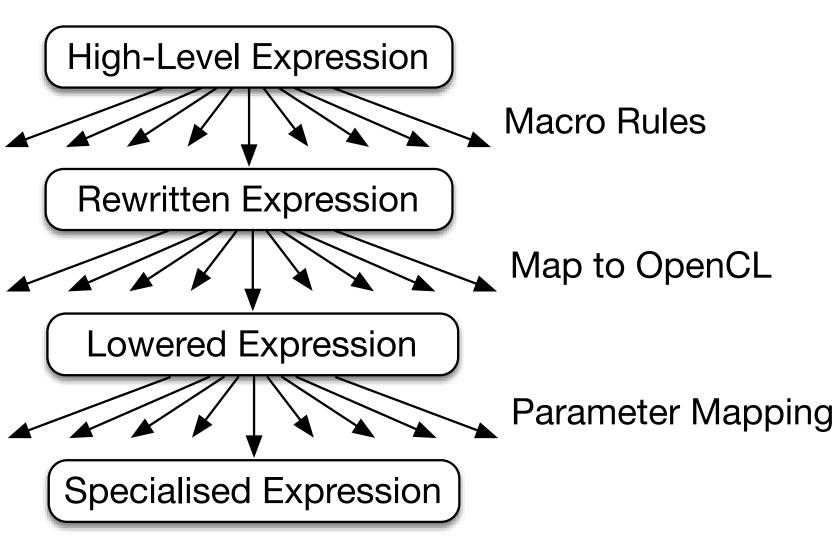


Register Blocking as a Rewrite Rules

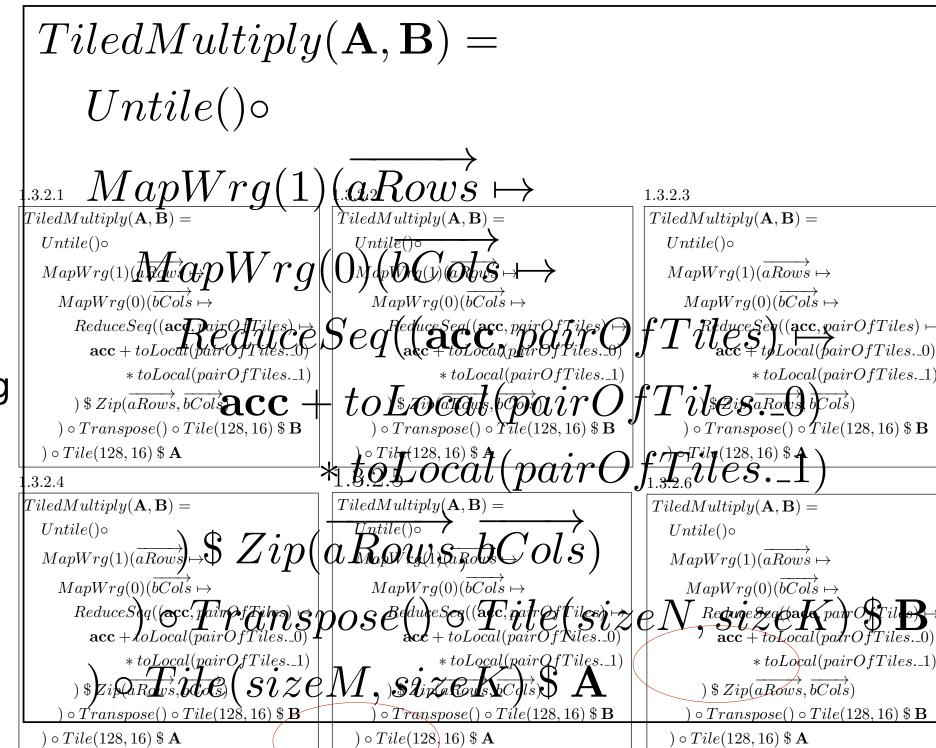
```
register Blocking =
    Map(f) \Rightarrow Join() \circ Map(Map(f)) \circ Split(k)
    Map(a \mapsto Map(b \mapsto f(a,b))) \Rightarrow Transpose() \circ Map(b \mapsto Map(a \mapsto f(a,b)))
    Map(f \circ g) \Rightarrow Map(f) \circ Map(g)
    Map(Reduce(f)) \Rightarrow Transpose() \circ Reduce((acc, x) \mapsto Map(f) \circ Zip(acc, x))
    Map(Map(f)) \Rightarrow Transpose() \circ Map(Map(f)) \circ Transpose()
    Transpose() \circ Transpose() \Rightarrow id
    Reduce(f) \circ Map(g) \Rightarrow Reduce((acc, x) \mapsto f(acc, g(x)))
    Map(f) \circ Map(g) \Rightarrow Map(f \diamond g)
                                              B
```

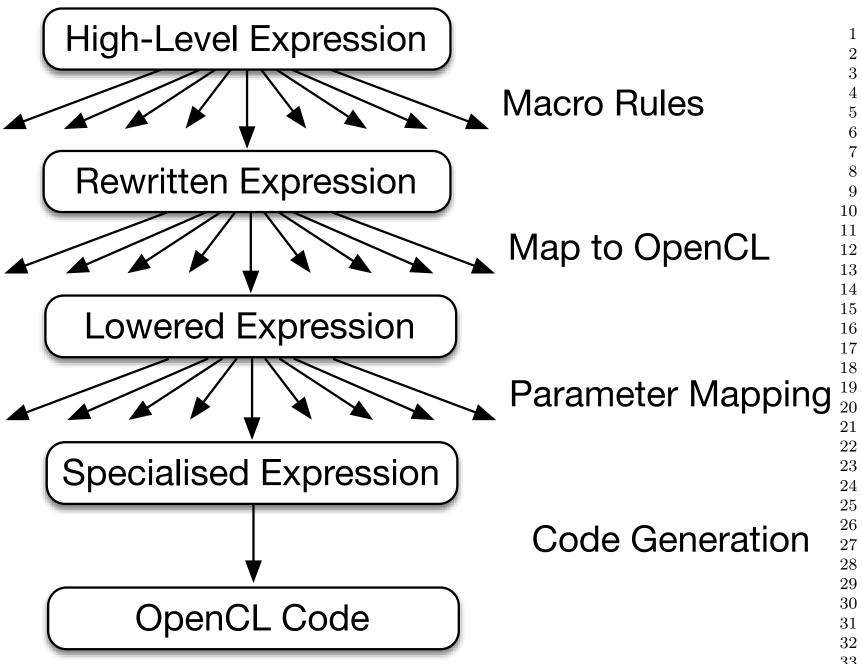






1.3.2





1.3.2.5

```
T_{	ext{local}}^{	ext{kernel mm_amd_opt(global float}} U_{	ext{l
private float acq_0; 1 ...; acc_31; private float blockOfA_0; ...; blockOfA_7;
           \begin{array}{l} \text{for (int } \mathbf{j} = \mathbf{0}; \ \mathbf{j} < \mathbf{8}; \ \mathbf{j} + \mathbf{+}) \ \{ \\ \text{blockOfA}_0 = \text{tileA}[\mathbf{0} + \mathbf{3} + \mathbf{3} + \mathbf{1} + \mathbf{1} + \mathbf{0}]; \\ \text{blockOfB}_0 = \text{tileB}[\mathbf{0} + \mathbf{0} + \mathbf{3} + \mathbf{1} + \mathbf{1} + \mathbf{0}]; \\ \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{blockOfB}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{blockOfB}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kOfA}_0 = \mathbf{0} \end{array} \\ \begin{array}{l} \text{of kOfA}_0 = \mathbf{0} \\ \text{of kO
                               ) Zip(aRows, bCols)
              C[48+8*lid1*N+64*w0+64*w1*N+0*N+lid0] = acc_3; ...; C[48+8*lid1*N+64*w0+64*w1*N+7*N+lid0] = acc_31
```

Heuristics for Matrix Multiplication

For Macro Rules:

- Nesting depth
- Distance of addition and multiplication
- Number of times rules are applied

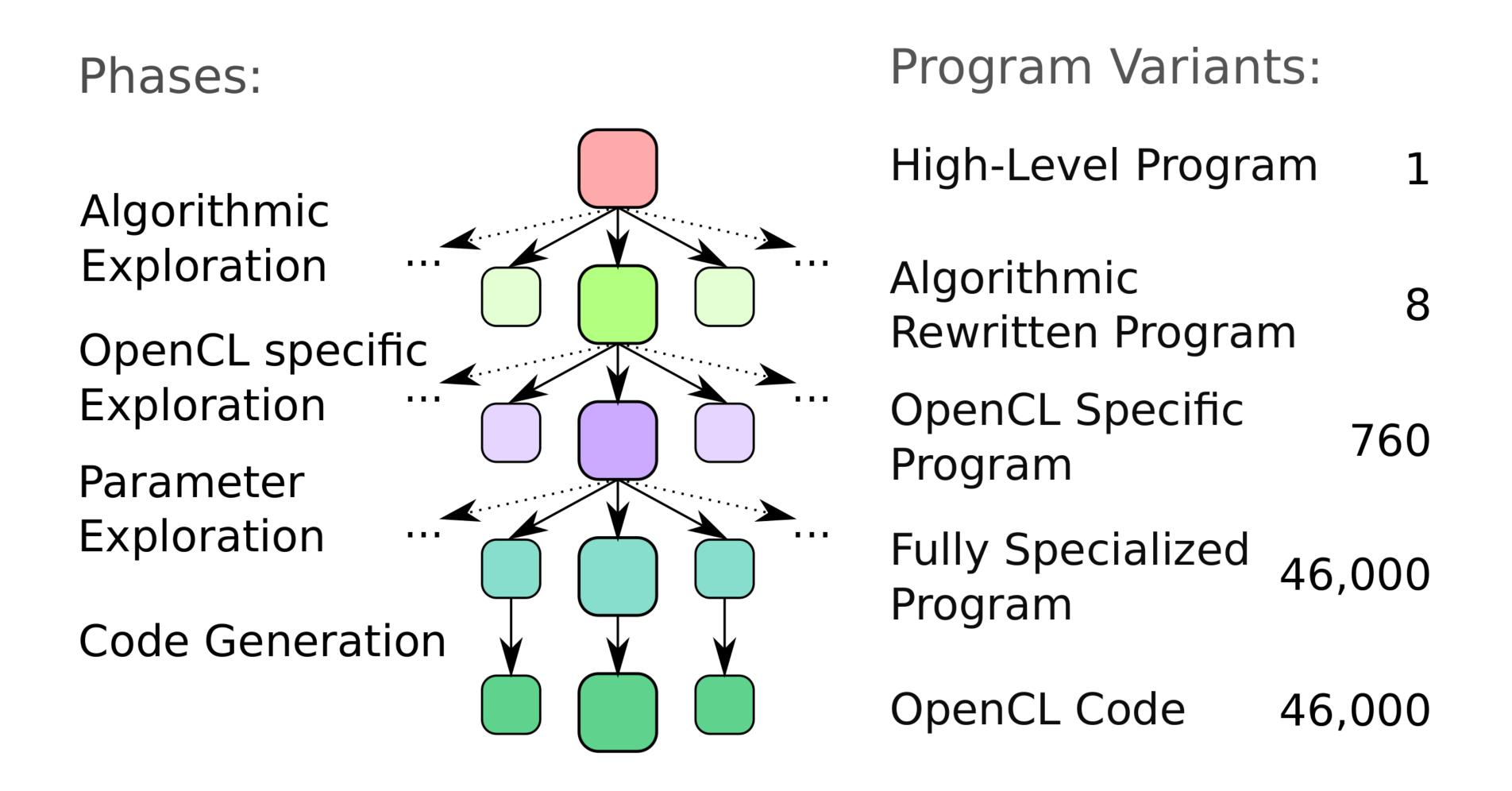
For Map to OpenCL:

- Fixed parallelism mapping
- Limited choices for mapping to local and global memory
- Follows best practice

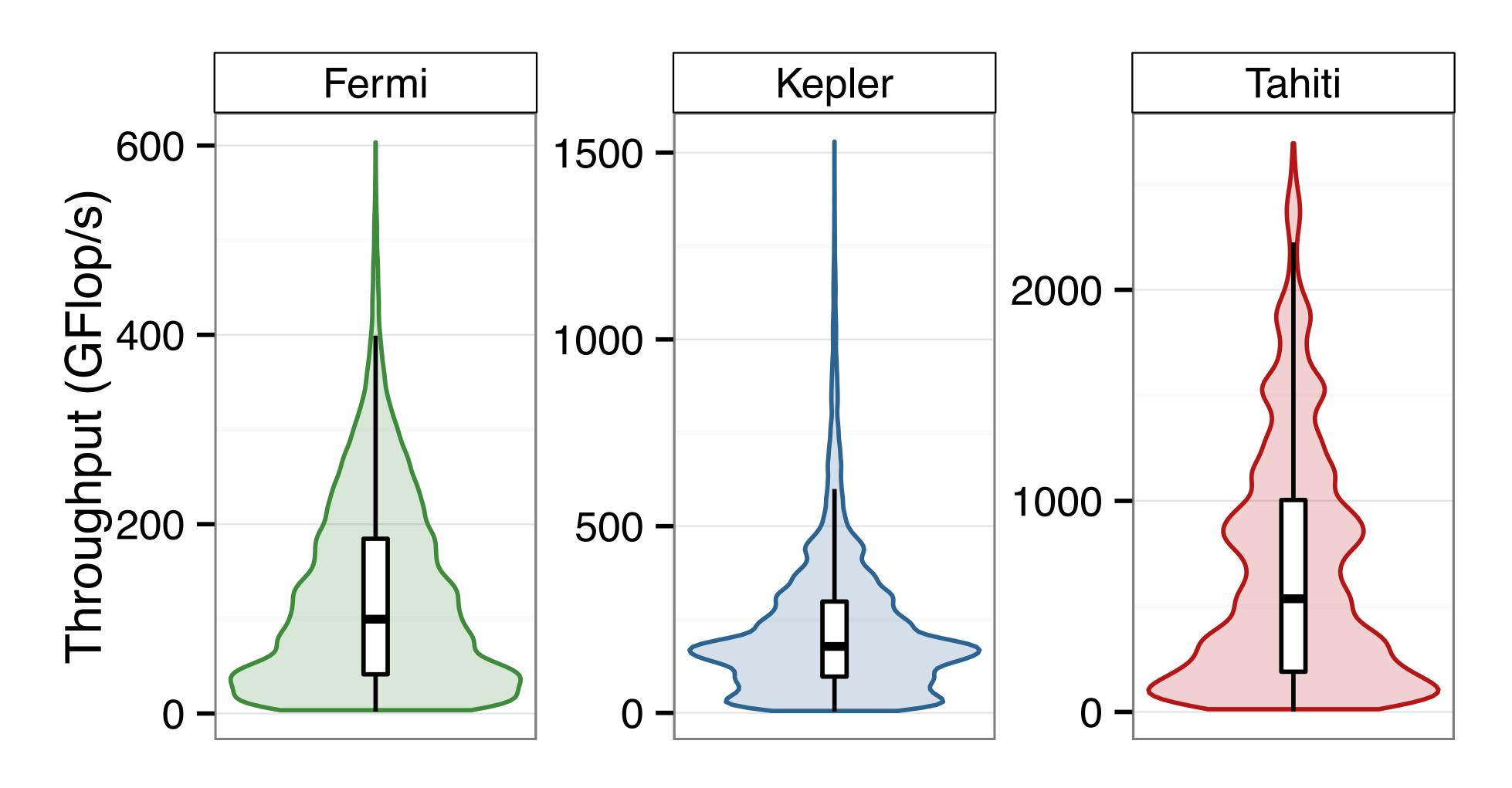
For Parameter Mapping:

- Amount of memory used
 - Global
 - Local
 - Registers
- Amount of parallelism
 - Work-items
 - Workgroup

Exploration in Numbers for Matrix Multiplication

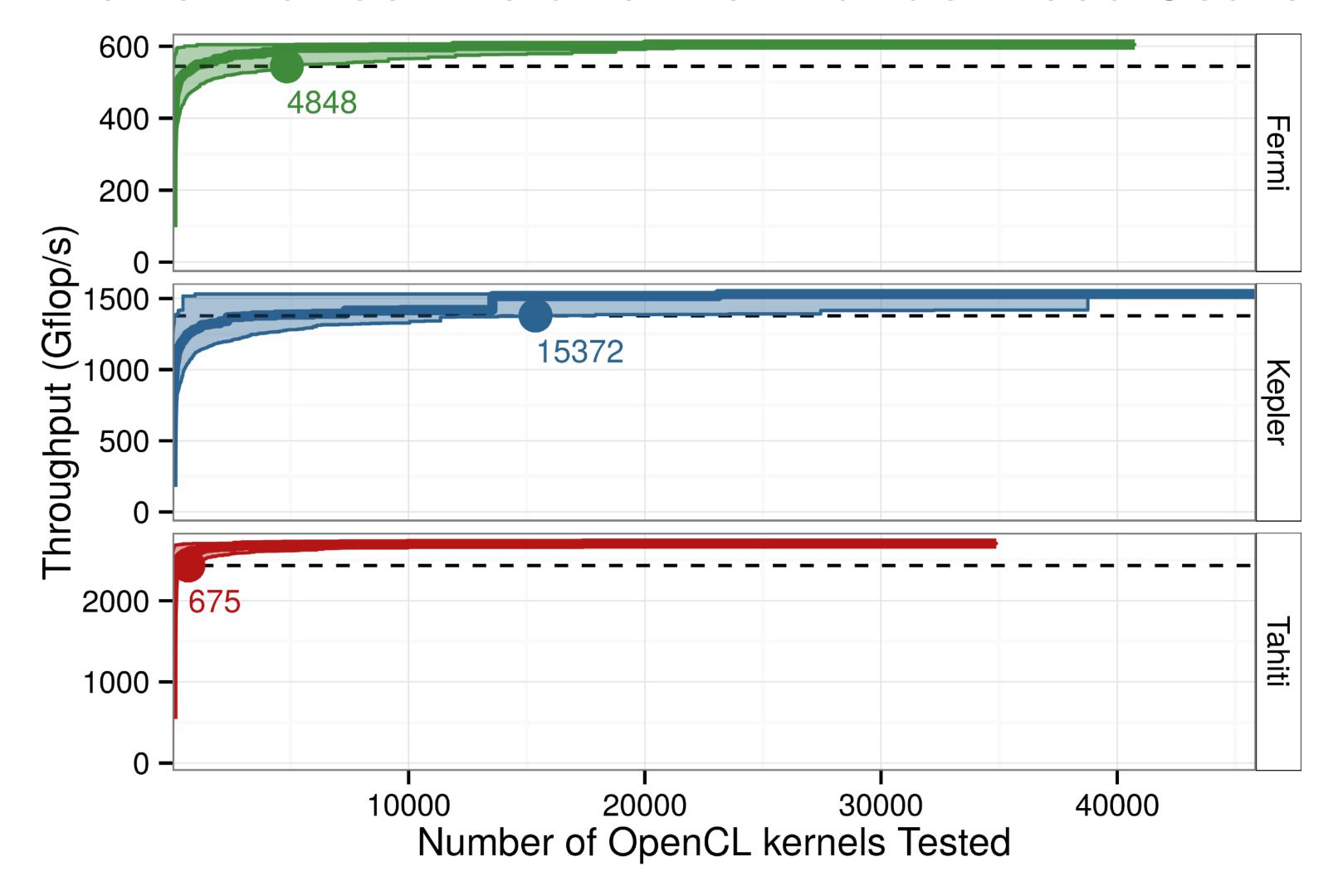


Exploration Space for Matrix Multiplication



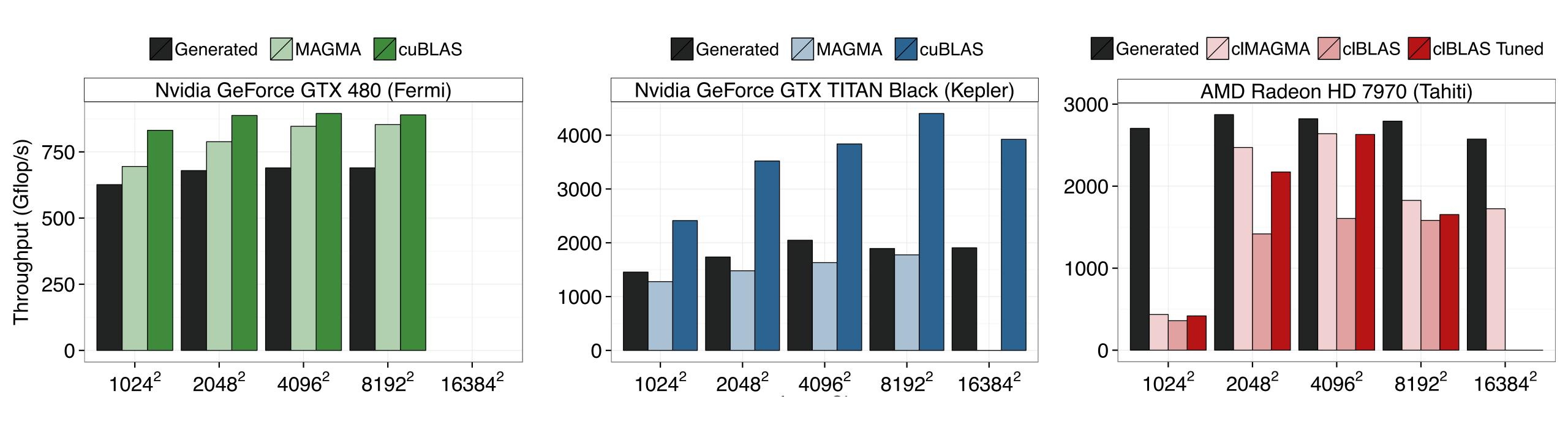
Only few OpenCL kernel with very good performance

Performance Evolution for Randomised Search



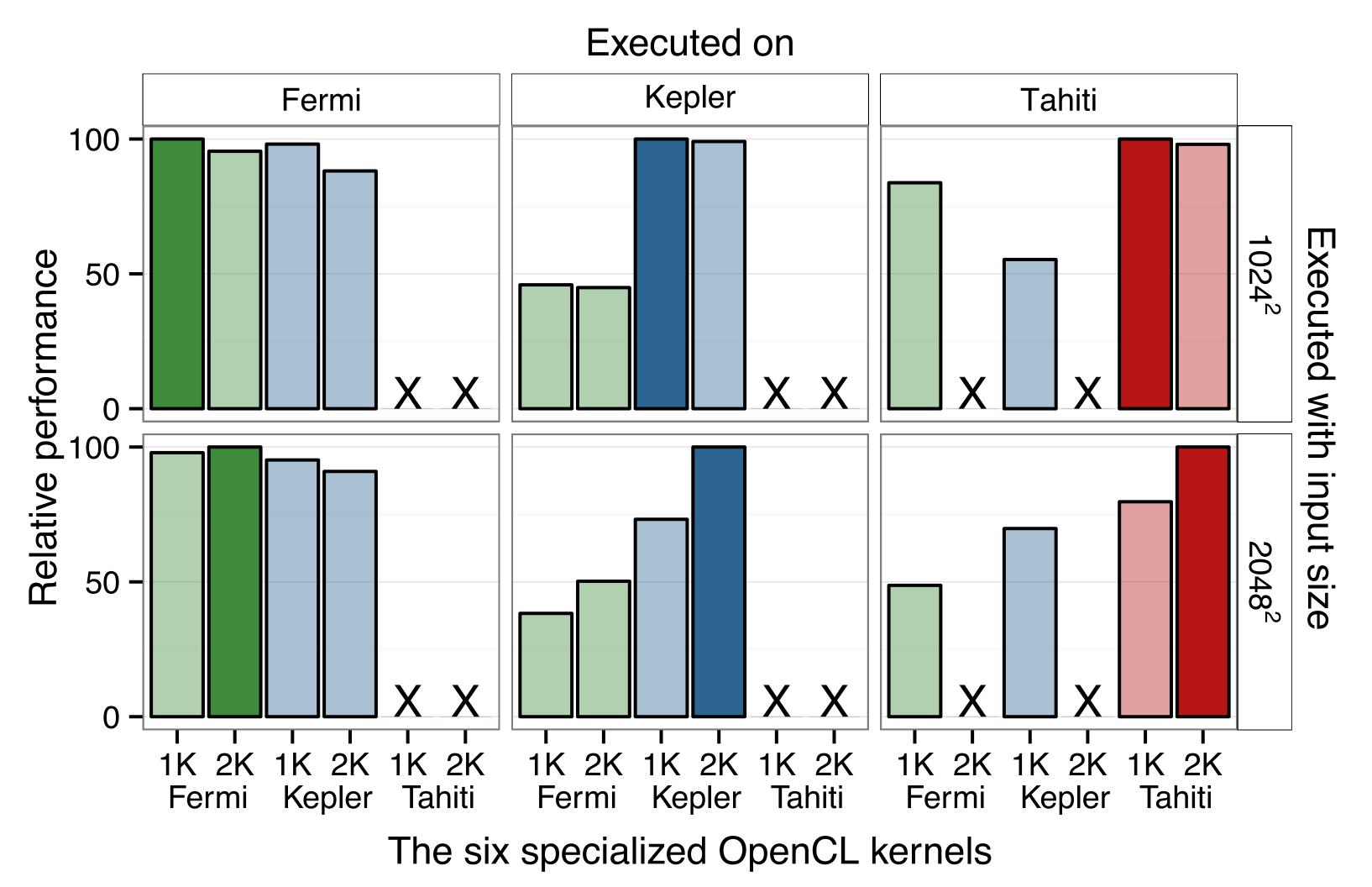
Even with a simple random search strategy one can expect to find a good performing kernel quickly

Performance Results Matrix Multiplication



Performance close or better than hand-tuned MAGMA library

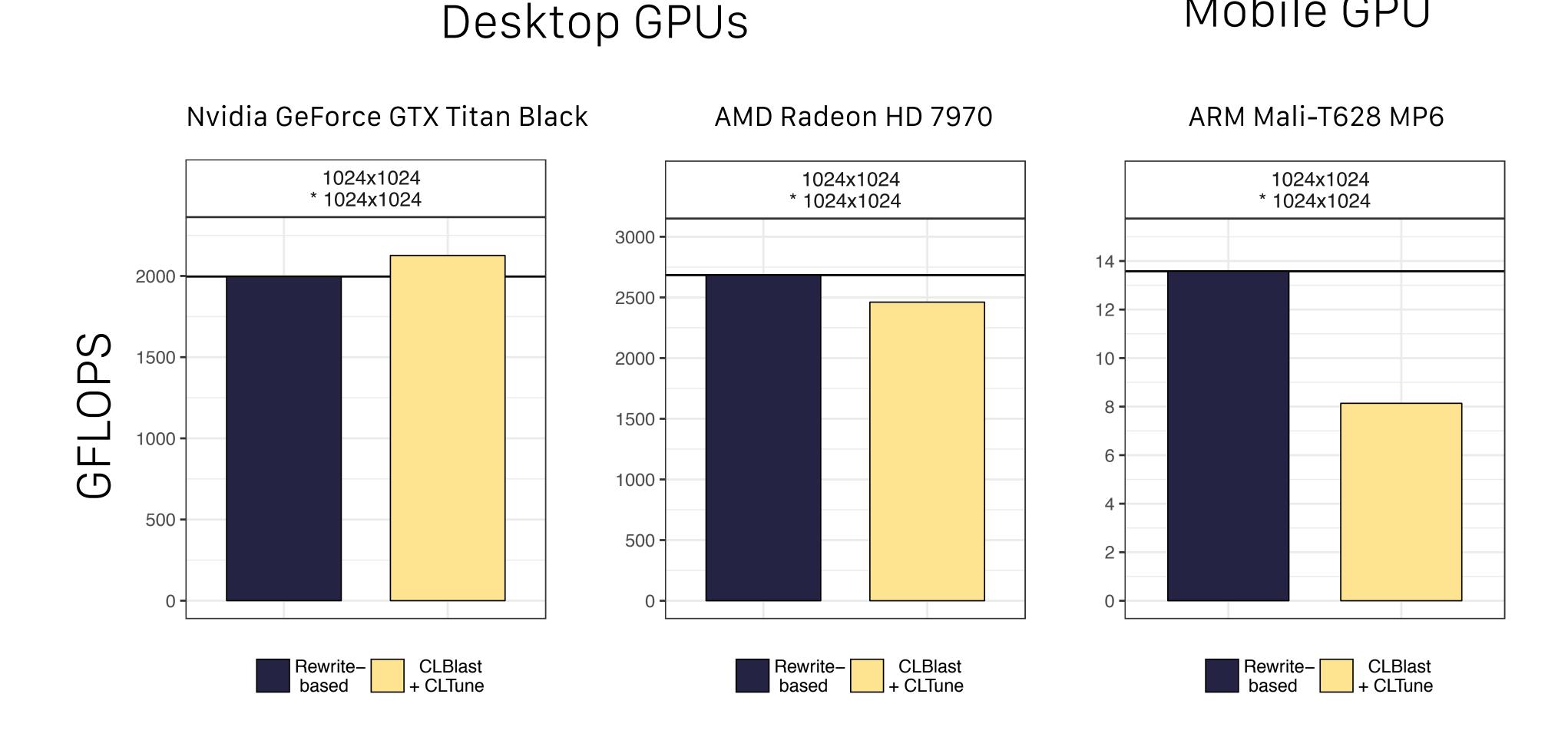
Performance Portability Matrix Multiplication



Generated kernels are specialised for device and input size

Desktop GPUs vs. Mobile GPU

Mobile GPU



Performance portable even for mobile GPU device!

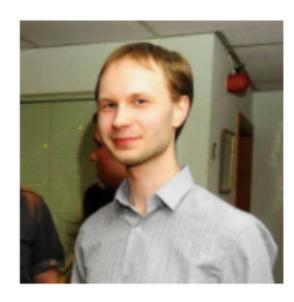




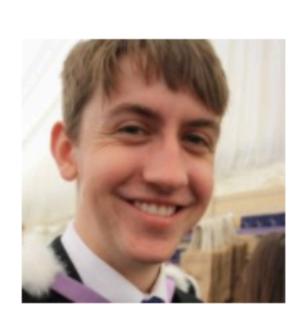




The LIFT Team













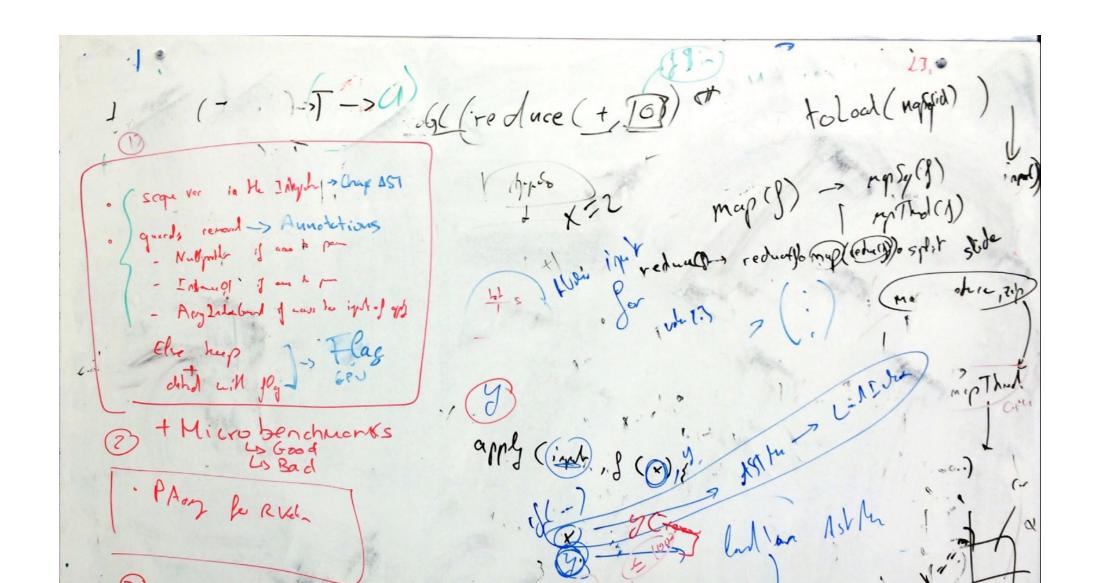


Federico Pizzuti PhD Student University of Edinburgh

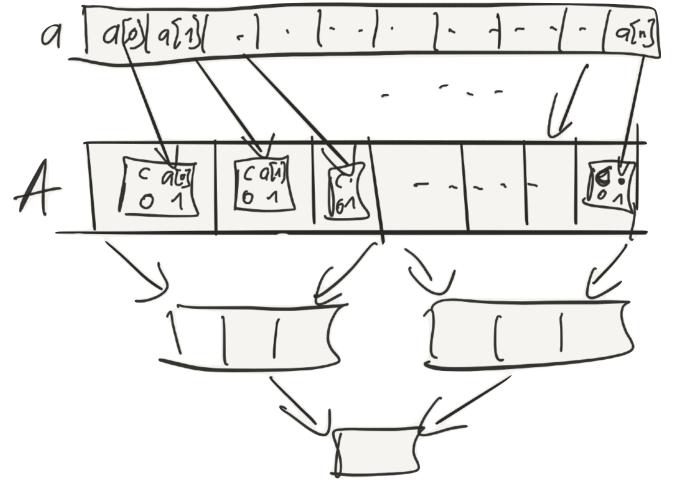
Identify hidden parallelism in LIFT programs

Parallelising non-associative reductions

Key idea: Rearrange data as matrices to exploit associative matrix multiplication

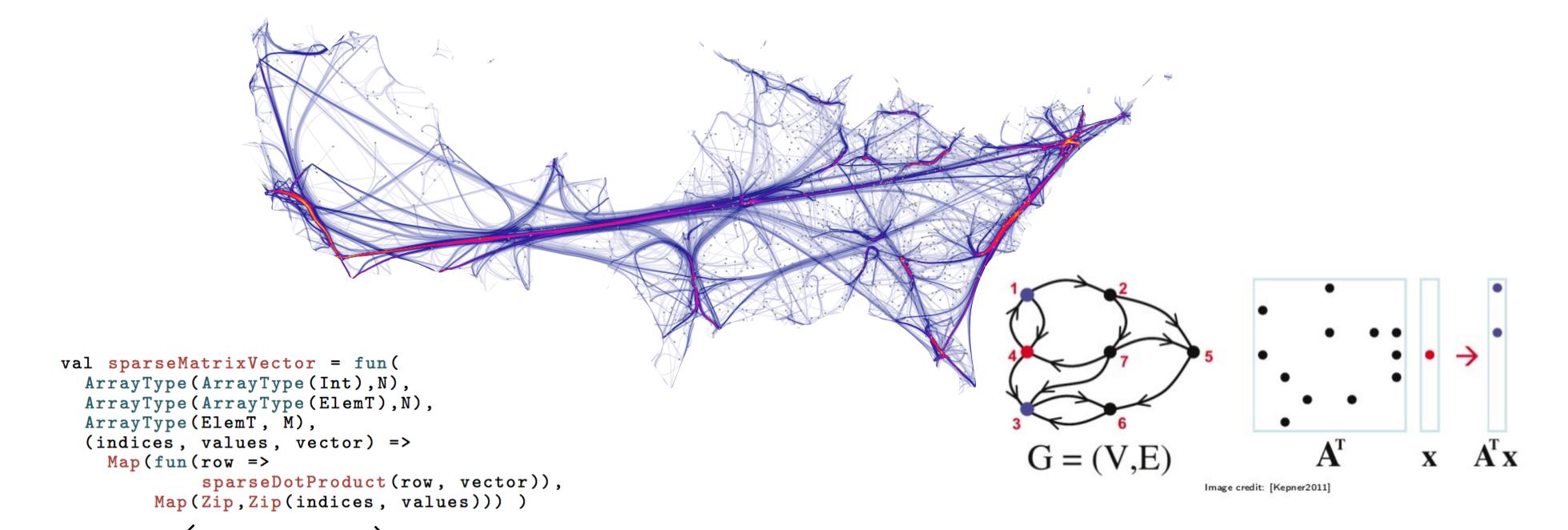


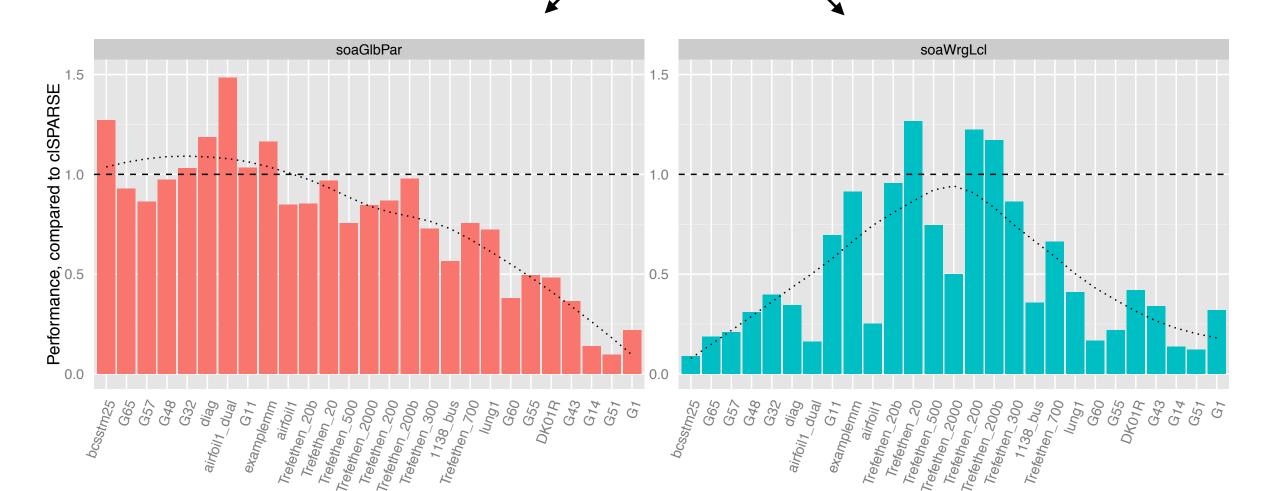
$$x \leftarrow 0$$
; for $i = 0$ to n do $x \leftarrow c \cdot x + a[i]$ done. $x \leftarrow x_0$; for $i = 0$ to n do $x \leftarrow A_i \times x$ done, where $x = \begin{pmatrix} x \\ 1 \end{pmatrix}$, $x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$, $A_i = \begin{pmatrix} c & a[i] \\ 0 & 1 \end{pmatrix}$.



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University of Edinburgh

Graph Algorithms via Sparse Linear Algebra in LIFT



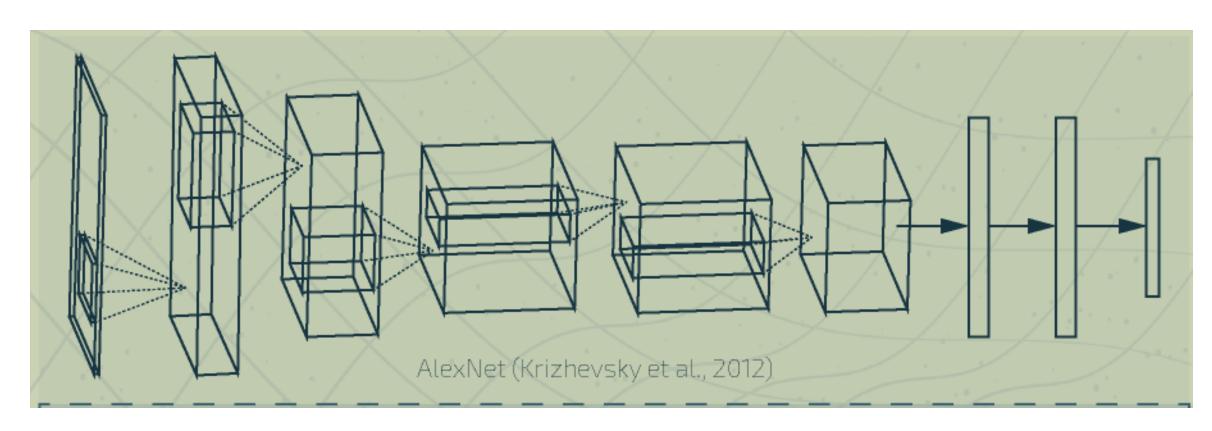


Differently optimised kernels for different inputs



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PhD Student
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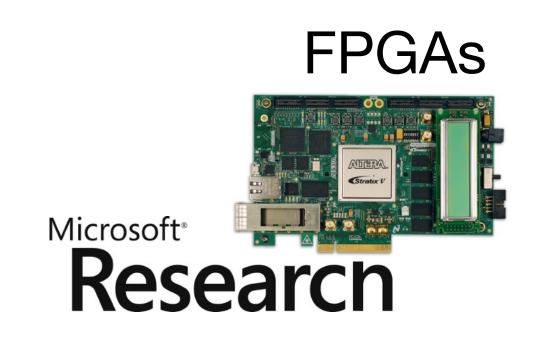
Optimising Deep Learning with LIFT

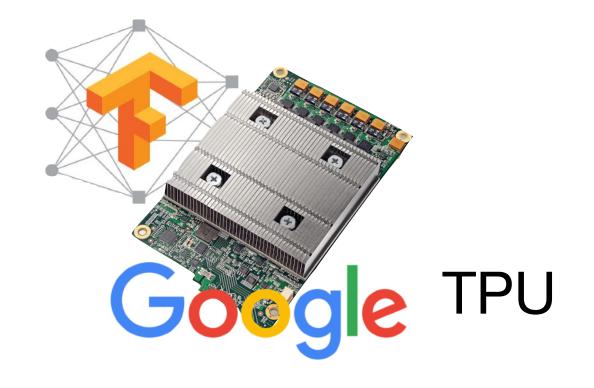


Express layers with LIFT primitives

```
fully_connected(f, weights, bias, inputs) :=
   Map((neuron_weights, neuron_bias) → f() o Reduce(add, neuron_bias) o
   Map(mult) $ Zip(inputs, neuron_weights)) $ Zip(weights, bias)
```

Optimise individual layers and across layers via rewrites











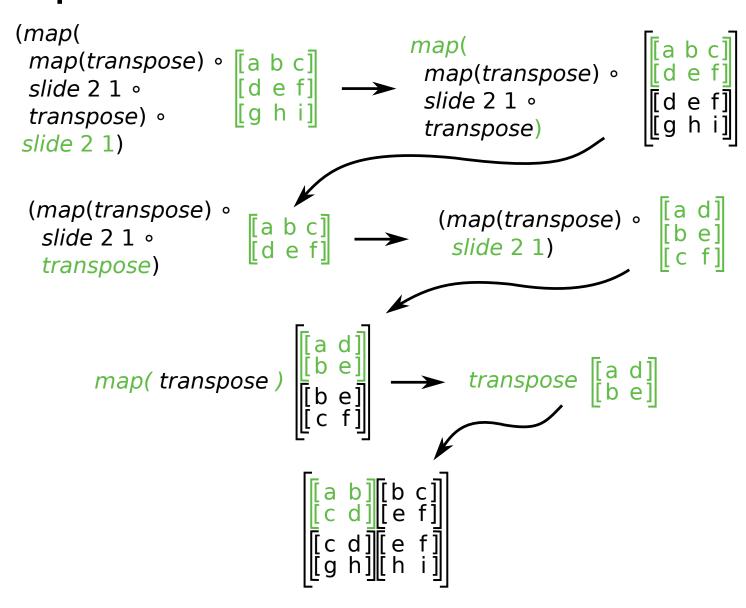
Bastian Hagedorn
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Larisa Stoltzfus
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University of Edinburgh

Stencil Computations in LIFT

Express Stencil with Skeletons



Explore optimisations as rewrites

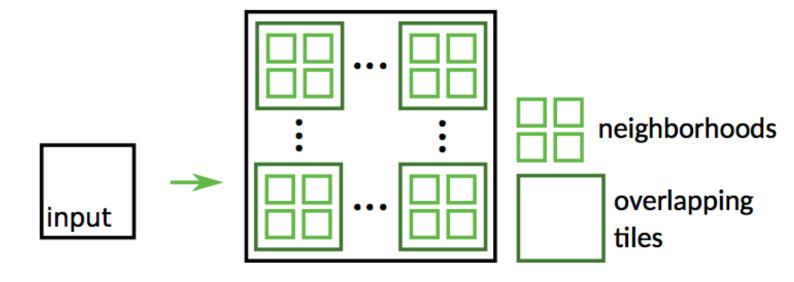
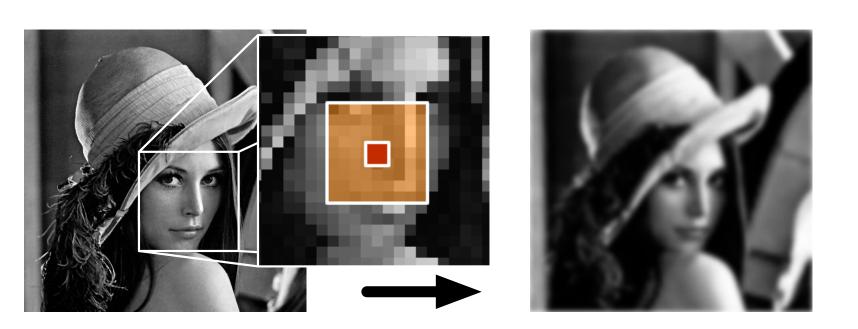
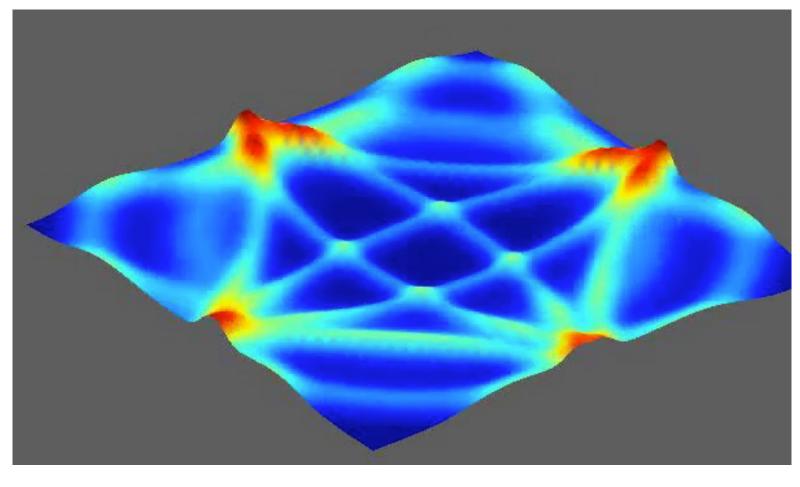


Image Processing



Acoustics Simulation



<u>Video</u>

Toomas Remmelg
PhD Student
University of Edinburgh

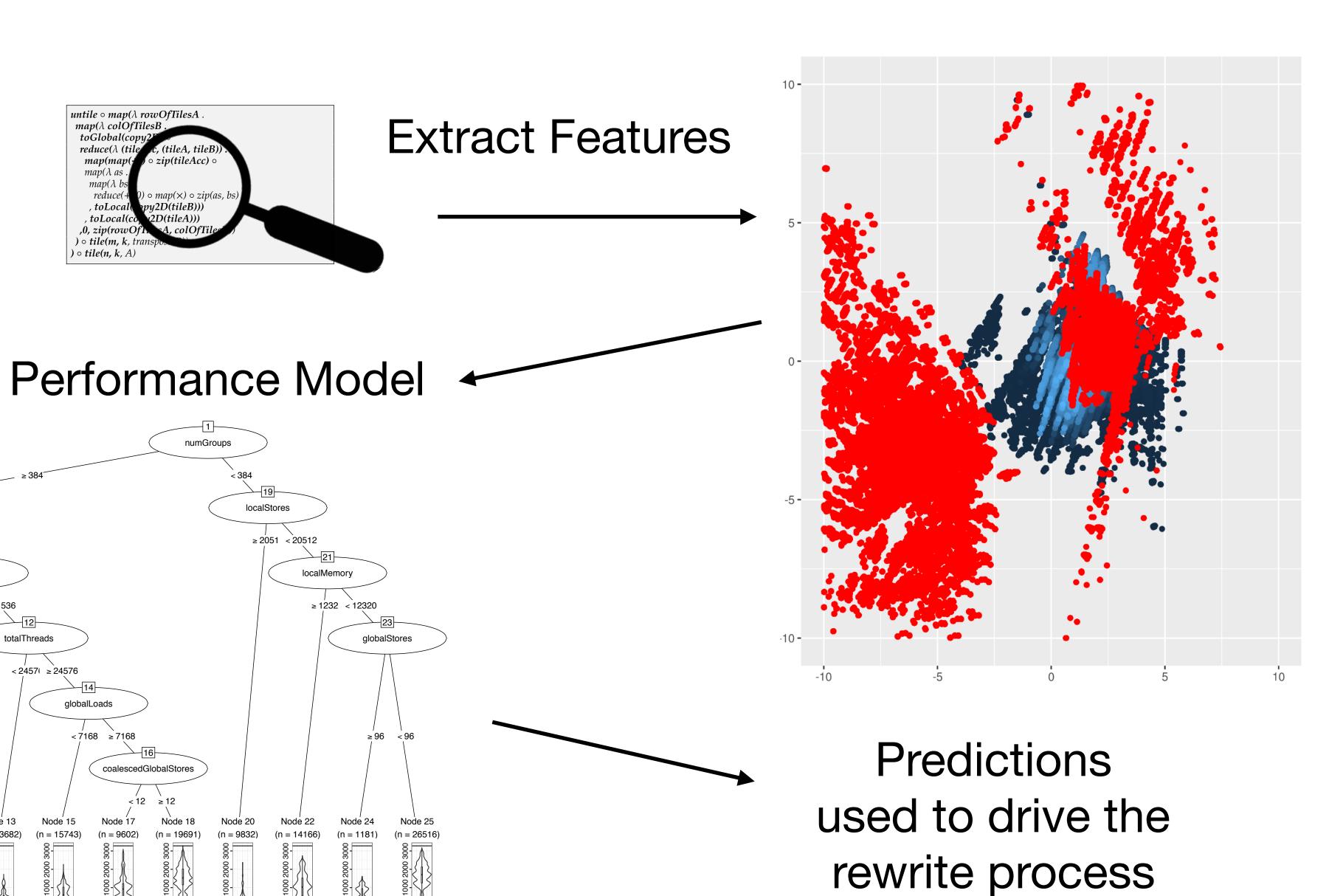
groupSize

totalLocalStores

≥ 132! < 1325400064

numGroups

Performance Modeling of LIFT Programs



Data Parallel Idealised Algol as new foundation for LIFT

Collaboration with Bob Atkey (Strathclyde) and Sam Lindley (Edinburgh)

- So far LIFT rewrites only functional expressions
- Data Parallel Idealised Algol (DPIA) combines functional and imperative constructs
- Allows formal translation of functional programs into efficient parallel imperative programs
- Types separate expressions, acceptors, and commands:
 - Expressions are purely functional computations
 - Acceptors describe modifiable locations in memory (≅ I-values in C)
 - Commands are imperative actions modifying memory

Draft Paper at: https://bentnib.org/dpia.html

```
reduce (+) 0 (map (\lambda x. fst x * \text{snd } x) (zip xs ys))
unctiona
     asScalar<sub>4</sub> (join (mapWorkgroup (\lambda zs_1. mapLocal (\lambda zs_2. reduce (\lambda x\ a. (fst x* snd x) + a) 0 (split 8192 zs_2)) zs_1)
                                        (split 8192 (zip (asVector<sub>4</sub> xs) (asVector<sub>4</sub> ys)))))
       parforWorkgroup (N/8192) (joinAcc (N/8192) 64 (asScalarAcc<sub>4</sub> (N/128) out)) (\lambda gid o.
         parforLocal 64 o (\lambda lid o.
           newPrivate num\langle 4 \rangle accum.
              accum.1 := 0;
Imperative
              for 2048 (\lambda i.
                accum.1 := accum.2 +
                   (fst (idx (idx (split 2048 (idx (split (8192 * 4) (zip (asVector<sub>4</sub> xs) (asVector<sub>4</sub> ys))) gid)) lid) i)) *
                   (snd (idx (split 2048 (idx (split (8192*4) (zip (asVector_4 xs) (asVector_4 ys))) gid)) lid) i)) );
              out := accum.2)
      kernel void KERNEL(global float *out, const global float *restrict xs,
                              const global float *restrict ys, int N) {
         for (int g_id = get_group_id(0); g_id < N / 8192; g_id += get_num_groups(0)) {
           for (int l_id = get_local_id(0); l_id < 64; l_id += get_local_size(0)) {</pre>
              float4 accum;
              accum = (float4)(0.0, 0.0, 0.0, 0.0);
              for (int i = 0; i < 2048; i += 1) {
                accum = (accum +
                            (vload4(((2048 * l_id) + (8192 * 4 * g_id) + i), xs) *
                             vload4(((2048 * l_id) + (8192 * 4 * g_id) + i), ys))); }
              vstore4(accum, ((64 * g_id) + l_id), out); } }
```

LIFT is Open-Source Software

http://www.lift-project.org/

https://github.com/lift-project/lift

